



Introduction to a Wait-Free Hash Table

Steven Feldman (UCF), Pierre LaBorde (UCF), Damian Dechev (Sandia, UCF)

Hash Table

- **Key-Value pair stores and retrieves data**
- **Three Core Functions:**
 - Put, Get, Delete
 - **Constant time: $O(1)$**

Design Goals:

- **Lock-Free and Wait-Free Nature**
 - All threads make progress
- **Concurrent/Efficient Resize**
- **Thread-Death Safety**
- **No CAS loops**

Lock-Free/Non-Blocking: ensure that no single thread can prevent another from accessing information.

Wait-Free: Guarantees that any process can complete any operation in a finite number of steps, regardless of the execution speeds of the other processes.

Features:

- Probing and non-probing versions
- **No global resize**
 - Concurrent expansion
- **Predictable memory usage**
 - Lookahead
- **Each key has a unique final position**



Brief Algorithm Description

Two types of Nodes:

- Spine Node: State, Memory Array
- Data Node: State, Key, Value
Key treated a bit string
Divided into blocks
Each block determines the position the key belongs in at each level.

Traversal Procedure:

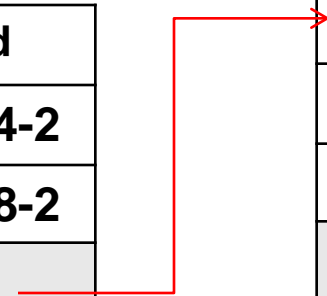
- 1) Bits in key are rearranged
- 2) First X bits determine the location
- 3) Examine the node
 - A. If Spine Node:
Using the next X bits, examine that spine's memory array
 - B. If Data Node:
Complete the thread's operation

Composed of Spine Nodes

- A position points to a spine or data node
- Max depth is X, where $2^X = \text{Spine Length}$

State: Old
0 K: 0-4-2
1 K: 1-8-2
2
3 K: 3-1-1
4 Spine
5 K: 5-5-5
6 NULL
7 K: 7-7-7

State: Old
0 K: 2-0-2
1 K: 2-1-2
2 NULL
3 NULL
4 K:2-4-4
5 NULL
6 NULL
7 K: 2-0-7

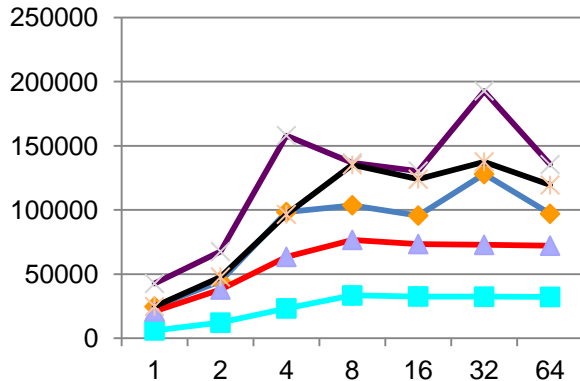




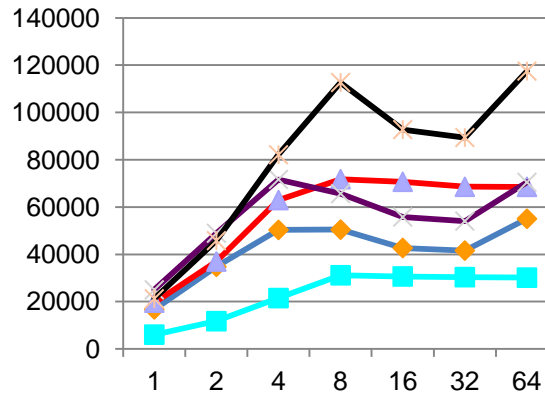
Performance Comparison

Operations Per Second

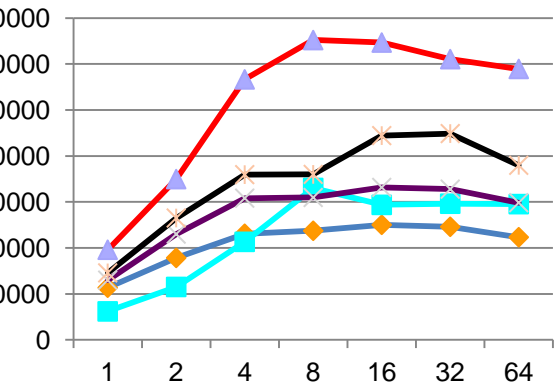
10% Modification



50% Modification



90% Modification



Number of Threads

Performance Analysis:

- Wait-Free performs better at Modification > 75%
- Algorithm still not optimized
- Probing version not worth it.

Legend:

- Wait-Free Hash Table (Non-Probing)
- Wait-Free Hash Table (Probing)
- TBB Concurrent Hashmap
- CDS Split List Hash Map (Lock-Free)
- Michaels Lock-Free Hash Map



Conclusion

Future Work:

- **Additional Functions**
 - Put-if-absent
 - Put-if-equals
 - Delete-if-equals
- **Memory pool**
- **Optimization of code**
- **Integration in real-world applications**

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Source Code: github.com/MrStevenFeldman/



Example

- Finding Key: 0-4-2
- 2-1-2 Hashes to 2-1-2
- Examines position 2
- Examines position 1
- Returns associated value

- Inserting Key: 1-2-3
- 1-2-3 Hashes to 5-6-1
- Position 5 is occupied
 - Adds and expands the table

