

Resource-aware Distributed Block-based LU Decomposition on Wireless



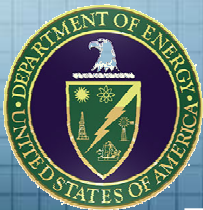
Sensor Networks



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Workshop

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(A) Approved for public release; distribution is unlimited.

Block-based LU

Any $N \times N$ matrix A can be partitioned into 4 matrices

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} L_{11} & 0 \\ L_{21} & L_{22} \end{bmatrix} \begin{bmatrix} U_{11} & U_{12} \\ 0 & U_{22} \end{bmatrix}$$

where A_{11} is $b \times b$ matrix

A_{12} is $b \times (N-b)$ matrix

A_{21} is $(N-b) \times b$ matrix

A_{22} is $(N-b) \times (N-b)$ matrix

Perform:

- LU decomposition (Crout's algorithm) over A_{11}

$$A_{11} = L_{11} \cdot U_{11}$$

- Upper decomposition over A_{12}

$$A_{12} = L_{11} \cdot U_{12}$$

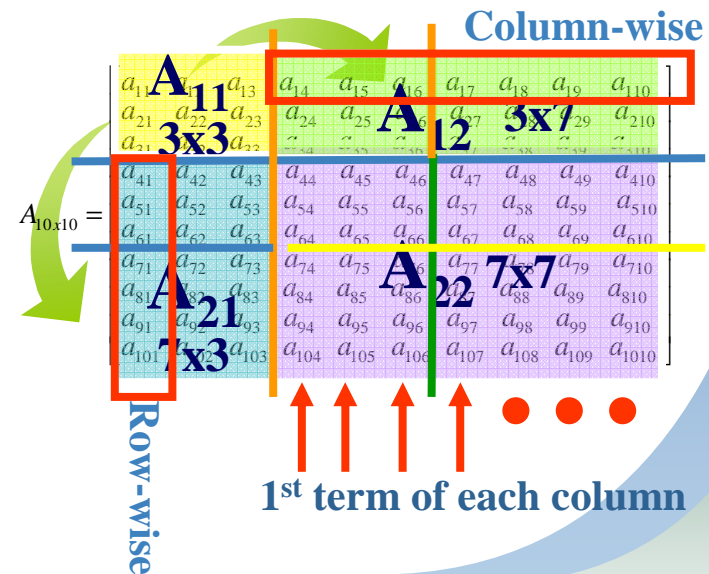
- Lower decomposition over A_{21}

$$A_{21} = L_{21} \cdot U_{11}$$

- Repeat the partitioning for A' or apply Crout's algorithm

$$A' = A_{22} - L_{21} \cdot U_{12}$$

For 10x10 Matrix A , and $b = 3$



Node N2
- responsible for **LU matrix operation**

Group1:
- Nodes responsible for upper matrix computation (**β operation**)
- Nodes have **odd addresses**

Group2
- Nodes responsible for lower matrix computation (**α operation**)
- Nodes have **even addresses**

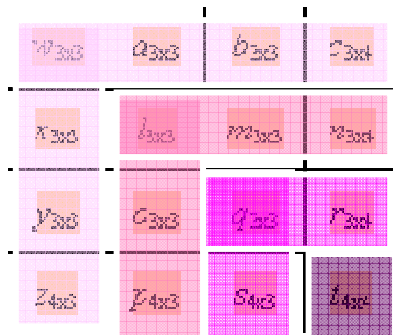
$$N = 13, b = 3$$

$W_{3 \times 3}$	$a_{3 \times 3}$	$b_{3 \times 3}$	$c_{3 \times 4}$
$X_{3 \times 3}$	$l_{3 \times 3}$	$m_{3 \times 3}$	$n_{3 \times 4}$
$Y_{3 \times 3}$	$u_{3 \times 3}$	$q_{3 \times 3}$	$r_{3 \times 4}$
$Z_{4 \times 3}$	$p_{4 \times 3}$	$s_{4 \times 3}$	$t_{4 \times 4}$

$$\#nodes = 2 \cdot \lfloor N/b \rfloor - 1 = 7$$



Round **3**

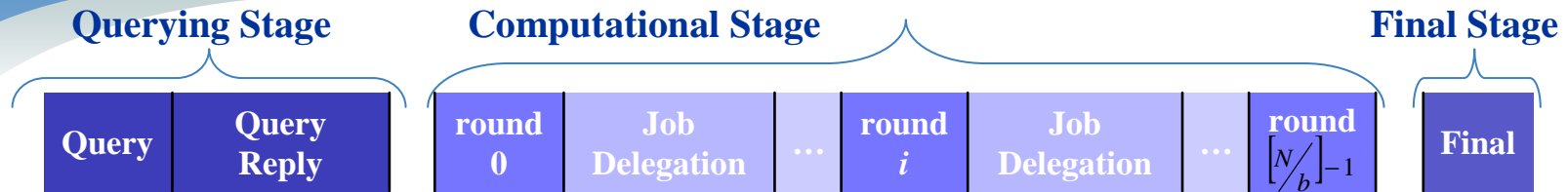


#Rounds = $\lceil N/b \rceil$
= 4

Round **i**

- BS transmits to grp1, grp2 and N2
- Each node in grp1 / grp2 broadcasts its result
- The 1st node in grp2 unicasts 1 message to N2
- At the end of each round, the 1st nodes in each group can **RETIRE**

N_2 $R_w = LU(w)$ RX: w TX: R_w	N_3 $R_a = \beta (a \text{ and } R_w)$ RX: a, R_w TX: R_a	N_5 $R_b = \beta (b \text{ and } R_w)$ $R_{5j} = R_a * R_b$ RX: b, R_a, R_a TX: R_a	N_7 $R_c = \beta (c \text{ and } R_w)$ $R_{7j} = R_a * R_c$ $R_{7j} = R_5 * R_c$ RX: c, R_w, R_a, R_p TX: R_a
N_4 $R_x = \alpha (x \text{ and } R_w)$ $R_{4j} = R_x * R_a$ RX: x, R_w, R_a TX: R_x, R_{4j}	N_2 $R_L = LU(L - R_{4j})$ RX: L, R_{4j} TX: R_L	N_5 $R_m = \beta [(m - R_{5j}) \text{ and } R_c]$ RX: m, R_L TX: R	N_7 $R_n = \beta [(n - R_{7j}) \text{ and } R_L]$ $R_{7j} += R_m * R_n$ RX: n, R_L, R_p TX: R
N_6 $R_y = \alpha (y \text{ and } R_n)$ $R_{6j} = R_y * R_a$ $R_{6j} = R_y * R_b$ RX: y, R_w, R_a, R_b TX: R_y	N_6 $R_o = \alpha [(o - R_{6j}) \text{ and } R_L]$ $R_{6j} += R_o * R_m$ RX: o, R_L, R_m TX: R_o, R_{6j}	N_3 $R_g = LU(g - R_{6j})$ RX: R_g TX: g, R_{6j}	N_7 $R_r = \mathcal{L} [(r - R_{7j}) \text{ and } R_g]$ RX: r, R_g TX: R_y
N_8 $R_z = \alpha (z \text{ and } R_w)$ $R_{8j} = R_z * R_a$ $R_{8j} = R_z * R_b$ $R_{8j} = R_z * R_c$ RX: z, R_w, R_a, R_b, R_c TX: R_z	N_8 $R_p = \alpha [(p - R_{8j}) \text{ and } R_L]$ $R_{8j} += R_p * R_m$ $R_{8j} += R_p * R_n$ RX: p, R_L, R_m, R_n TX: R_p	N_3 $R_s = \alpha [(s - R_{8j}) \text{ and } R_g]$ $R_{8j} += R_s * R_r$ RX: s, R_g, R_r TX: R_s, R_{8j}	N_2 $R_t = LU(t - R_{8j})$ RX: t, R_{8j} TX: R_a



1. Querying stage:

- BS broadcasts QUERY packet to query the **voltage level** and the amount of **available memory**
- BS selects $2 \cdot \lfloor N/b \rfloor - 1$ nodes to be **active** and the rest to **standby** (used in a failure management scheme through auto-power notification)

2. **Computation Stage:** It is formed of $\lfloor N/b \rfloor$ rounds and $\lfloor N/b \rfloor - 1$ job delegation periods

3. **Final Stage:** BS regroups the result

Each node infers its time slot from its ID

➤ Energy and Storage-aware Job Allocation:

Divide voltage levels 3 levels

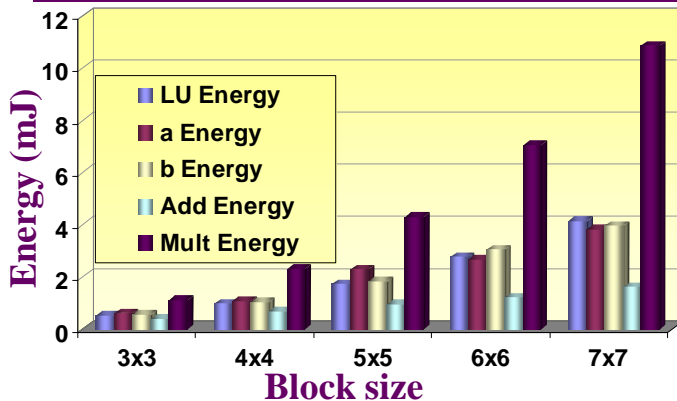
- $10\% < V1 = V_{min} < 30\%$ of the battery voltage
- $V_{min} < V2 < V_{max}$ of the battery voltage
- $V3 = V_{max} > 70\%$ of the battery voltage

Divide amount of free memory into 3 levels

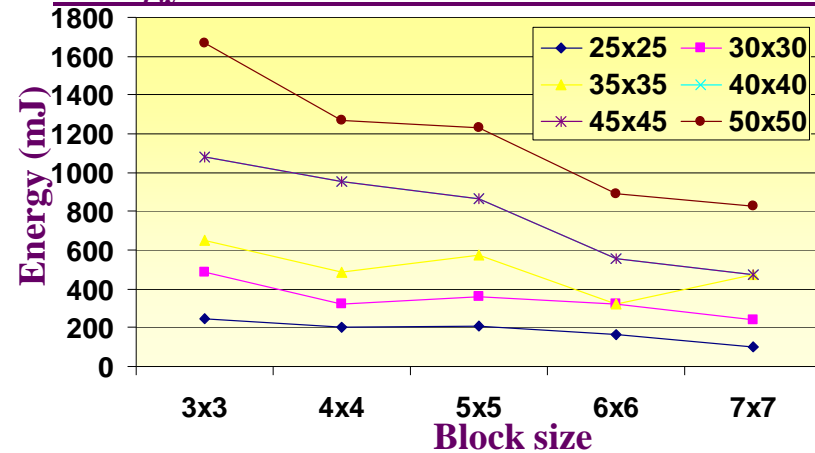
- $10\% < M1 = M_{min} < 30\%$ of memory
- $M_{min} < M2 < M_{max}$ of memory
- $M3 = M_{max} > 70\%$ of memory

ENERGY and STORAGE requirements **increase** with the iteration number

Energy per Operation vs. Matrix Size



Energy for diff. matrix sizes and block sizes



$P_{avg} \approx 4.06 \times 10^{-2} W$

$E_{max} = 10.87 \text{ mJ}, E_{min} = 0.51 \text{ mJ}$

P_{max} ranged between $4.09 \times 10^{-2} W$ and $4.12 \times 10^{-2} W$

Contributions

- ✓ Proposed a scheme for distributing block-based LU decomposition on WSN
- ✓ Proposed an energy and resource-aware job allocation
- ✓ Proposed a mostly-decentralized synchronous MAC / Time Slot Allocation
- ✓ Proposed a failure management scheme through auto-power notification
- ✓ Proved that LU decomposition requirements are feasible on WSN