

# High-Productivity Stream Programming for High-Performance Systems

Rodric Rabbah, Bill Thies, Michael Gordon, Janis Sermulins,  
and Saman Amarasinghe

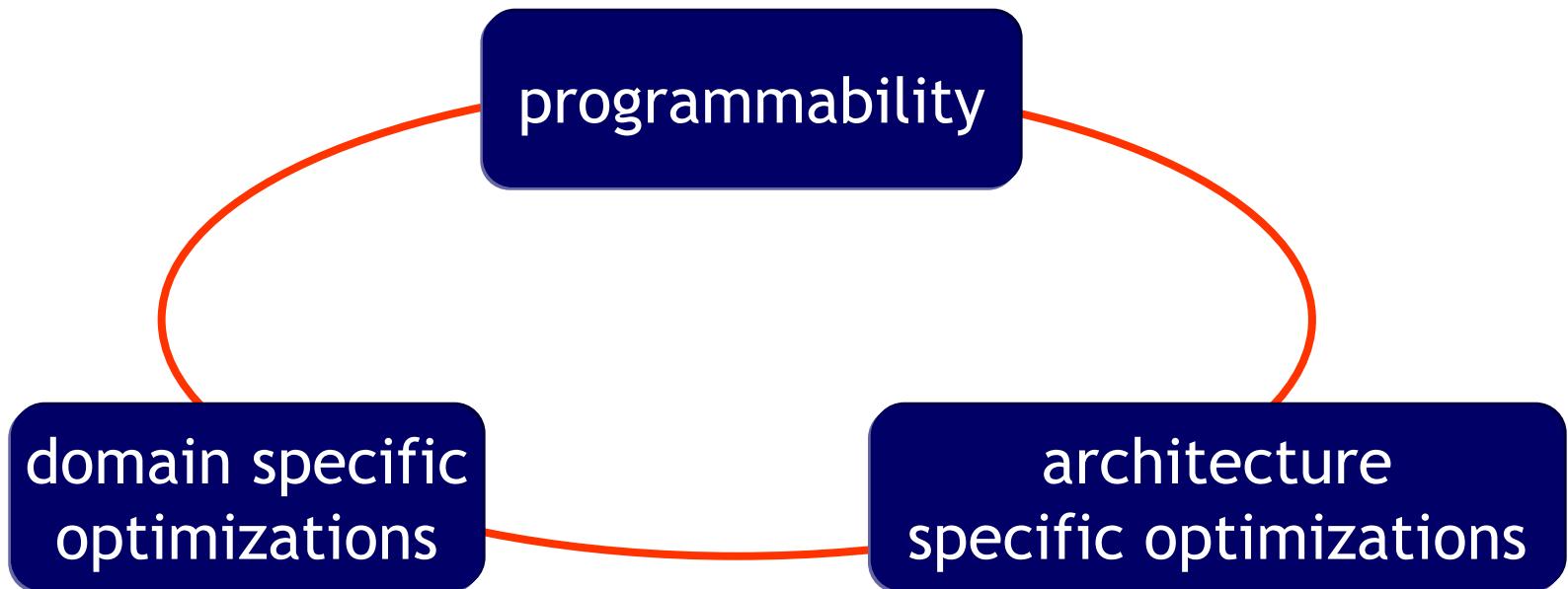
Massachusetts Institute of Technology



HPEC 2005, MIT LL

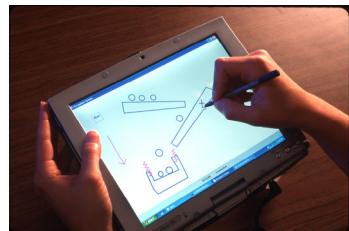
# The StreamIt Vision

- Boost productivity, enable faster development and rapid prototyping



- Simple and effective optimizations for streams
- Targeting tiled architectures, clusters of workstations, DSPs, and traditional uniprocessors

# Why an Emphasis on Streaming?

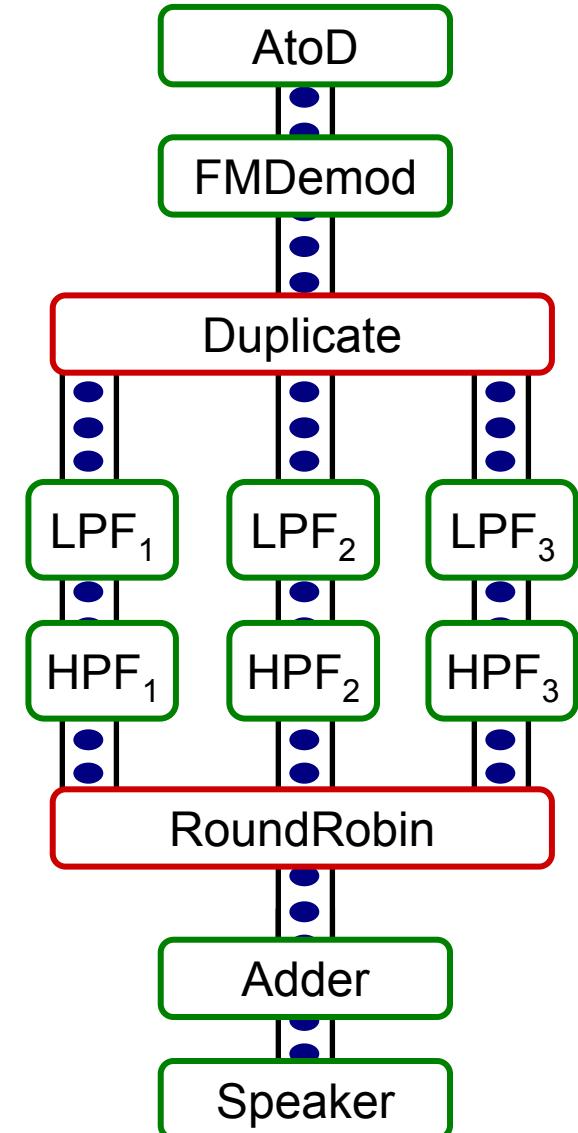


# Streaming in other Domains as well

- Cryptography
- Databases
- Face recognition
- Network processing and security
- Scientific codes
- ...
- Attractive programming model because of a simple mapping from specification to implementation

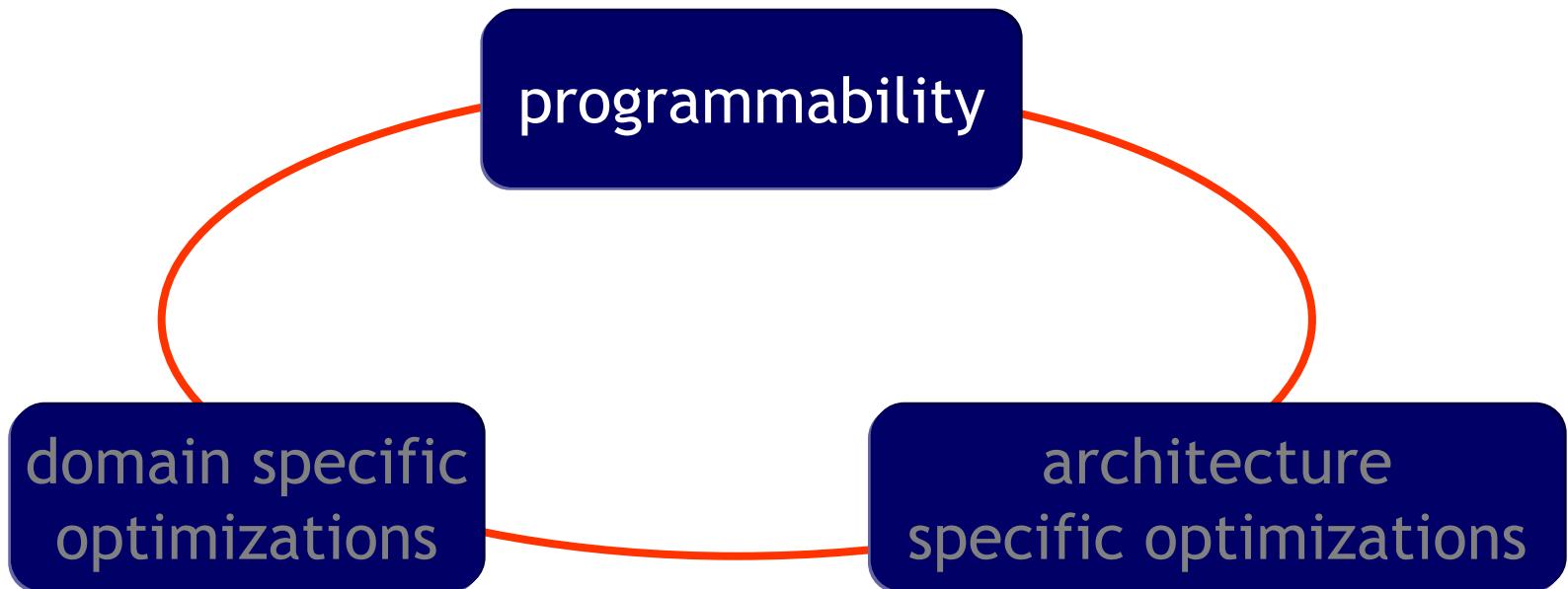
# Properties of Stream Programs

- Mostly regular and repeating computation
- Parallel, independent computation with explicit communication
- Amenable to aggressive compiler optimizations  
[ASPLOS '02, PLDI '03, LCTES'03, LCTES '05]



# The StreamIt Vision

- Boost productivity, enable faster development and rapid prototyping



- Simple and effective optimizations for streams
- Targeting tiled architectures, clusters of workstations, DSPs, and traditional uniprocessors

# Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Natural correspondence  
between text and  
application graph

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
}
```

```
    join roundrobin();
```

```
}
```

```
    add Adder();
```

```
    add Speaker();
```

```
}
```

# Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Streams are easily composed

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
}
```

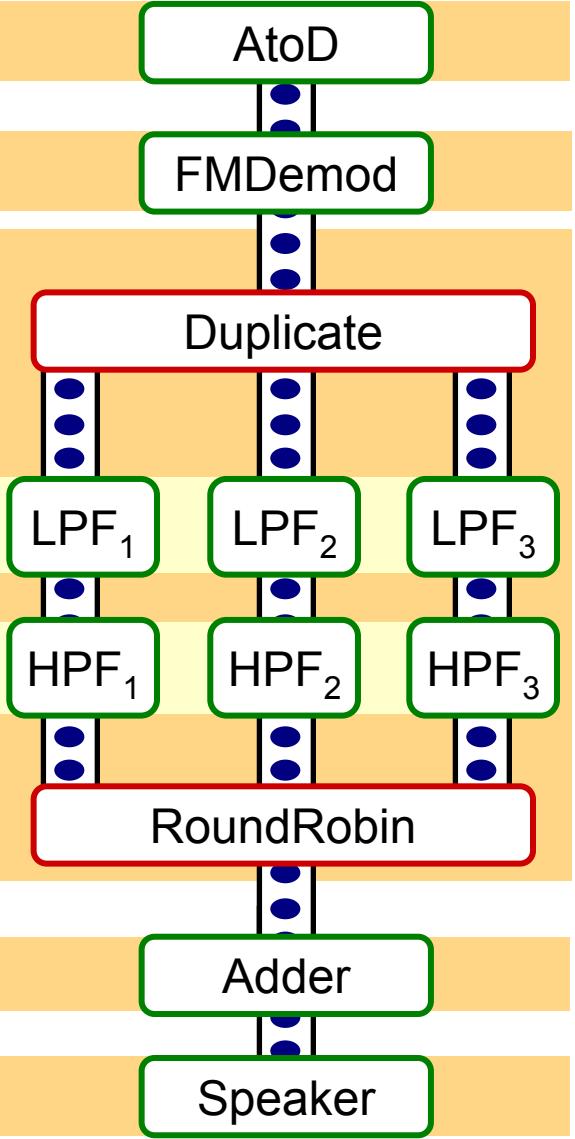
```
}
```

```
    join roundrobin();
```

```
}
```

```
    add Adder();
```

```
    add Speaker();
```



```
}
```

# Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Streams are parameterized, and malleable

```
AtoD
↓
FMDemod
↓
Duplicate
↓
LPF1   LPF2   LPF3
↓       ↓       ↓
HPF1   HPF2   HPF3
↓       ↓       ↓
RoundRobin
↓
Adder
↓
Speaker
```

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
}
```

```
}
```

```
}
```

```
    add Adder();
```

```
}
```

# Programming in StreamIt

```
void->void pipeline FMRadio(int N, float freq1, float freq2) {
```

```
    add AtoD();
```

- Application is architecture independent (i.e., portable)

```
    add FMDemod();
```

```
    add splitjoin {
        split duplicate;
        for (int i=0; i<N; i++) {
            add pipeline {
```

```
                add LowPassFilter(freq1 + i*(freq2-freq1)/N);
```

```
                add HighPassFilter(freq2 + i*(freq2-freq1)/N);
```

```
            }
```

```
        }
```

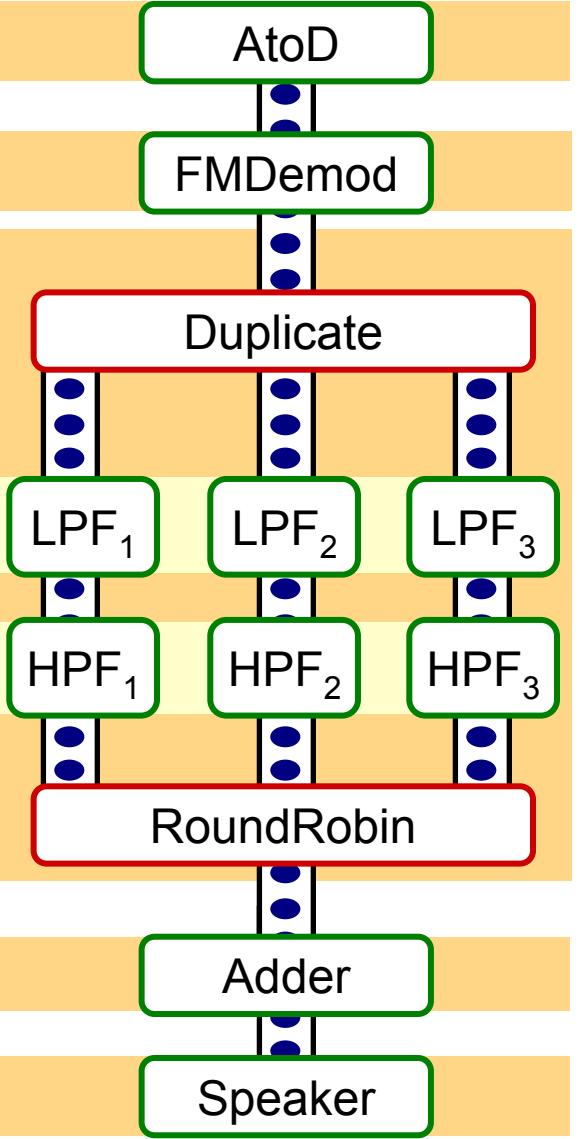
```
        join roundrobin();
```

```
}
```

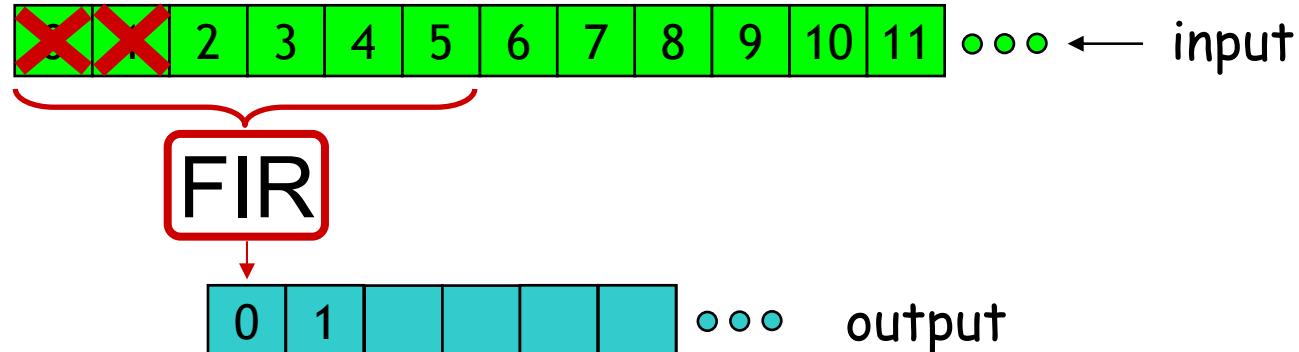
```
    add Adder();
```

```
    add Speaker();
```

```
}
```



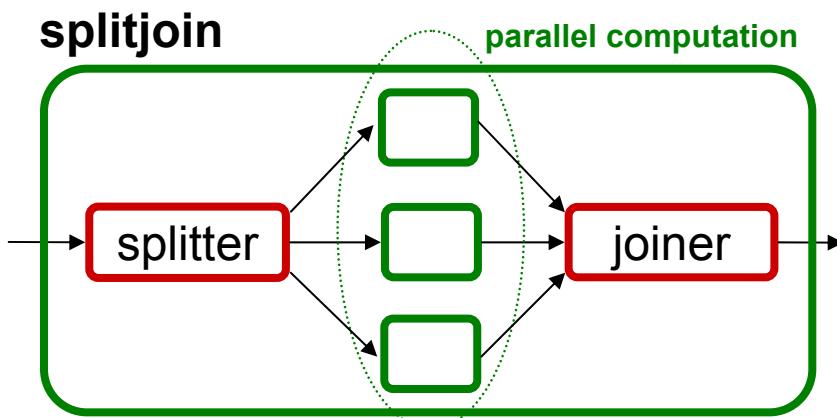
# Filters as Computational Elements



```
float→float filter FIR (int N) {
    work push 1 pop 1 peek N {
        float result = 0;
        for (int i = 0; i < N; i++) {
            result += weights[i] * peek(i);
        }
        push(result);
        pop();
    }
}
```

# Benefits of StreamIt

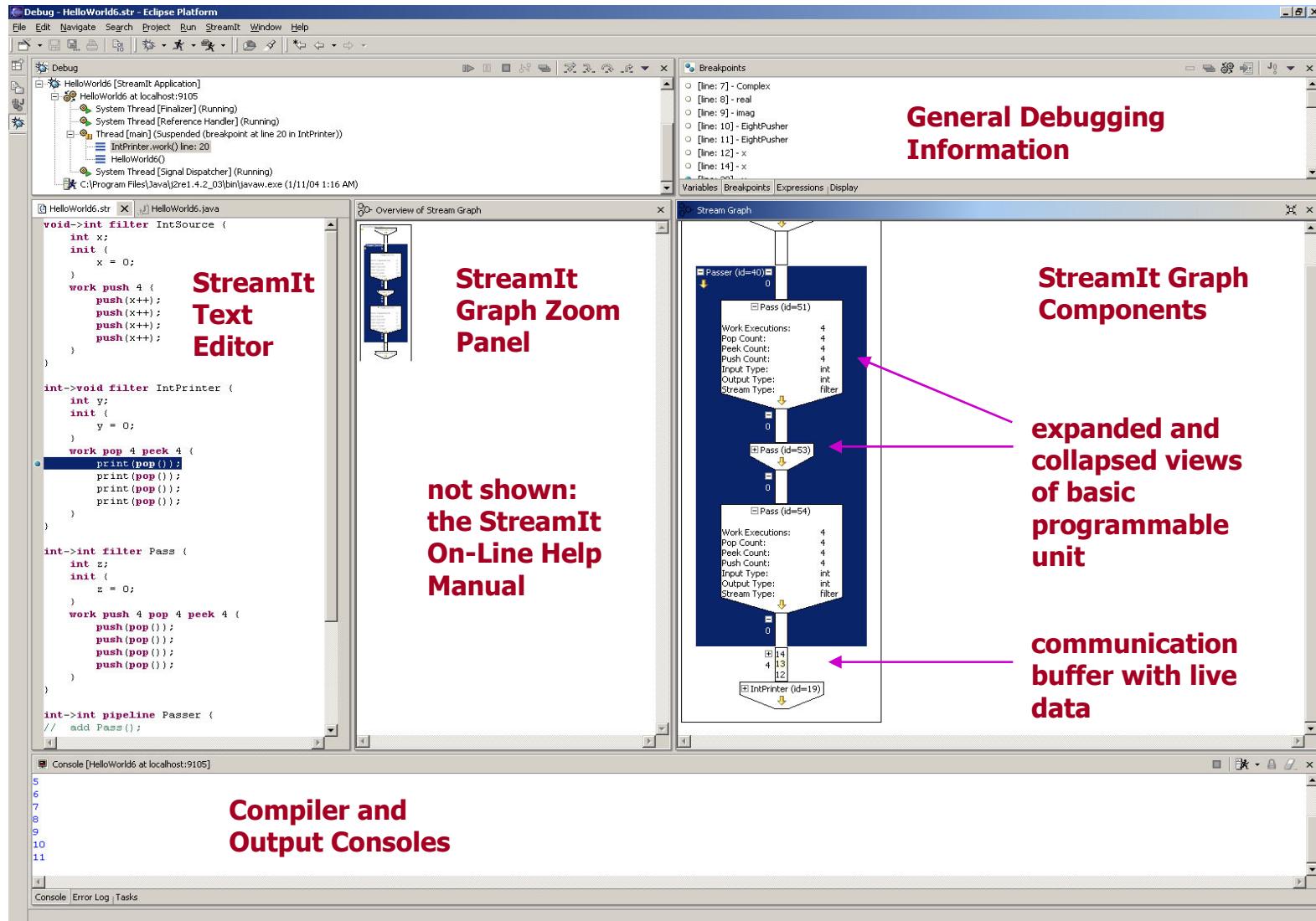
- Communication is exposed and pipeline parallelism is more readily discovered
- Flow of data provides a frame of reference for reasoning about “time” [PPoPP ’05]
  - Powerful advantage when debugging parallel programs



versus

- Multiple threads with independent program counters
- Non-deterministic execution

# StreamIt Development Environment



General Debugging Information

StreamIt Graph Components

expanded and collapsed views of basic programmable unit

communication buffer with live data

not shown:  
the StreamIt  
On-Line Help  
Manual

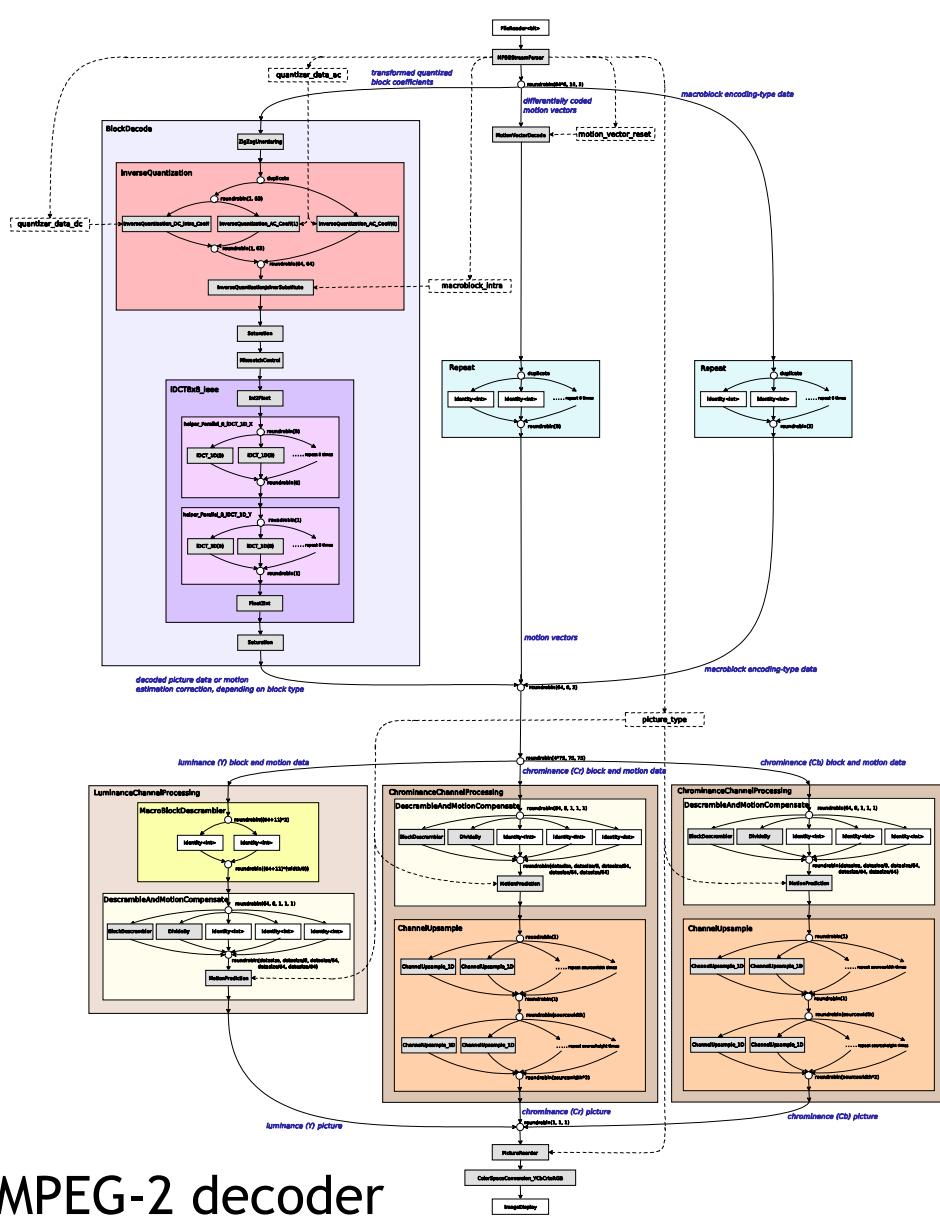
Compiler and  
Output Consoles

[PHEC '05]

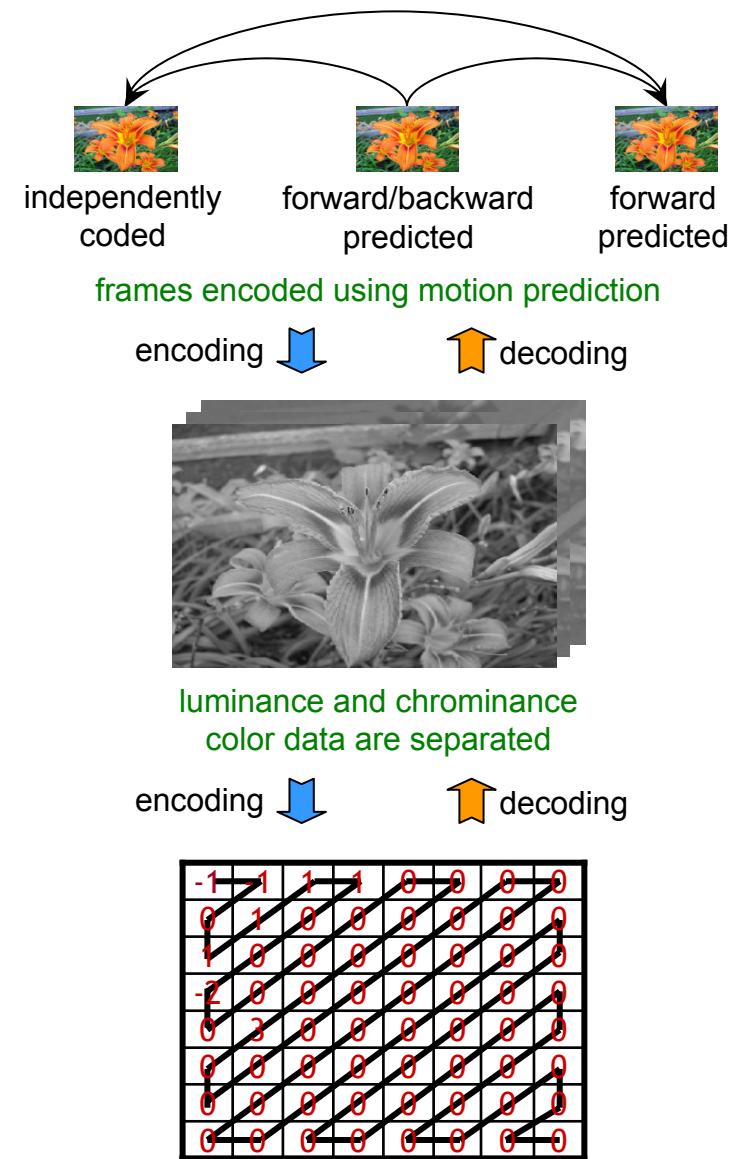
# StreamIt Applications

- Software radio
- Frequency hopping radio
- Acoustic beam former
- Vocoder
- GMTI (ground moving target indicator)
- DES and Serpent blocked ciphers
- Sorting
- FFTs and DCTs
- JPEG
- ...

# MPEG: Motion Video Codec

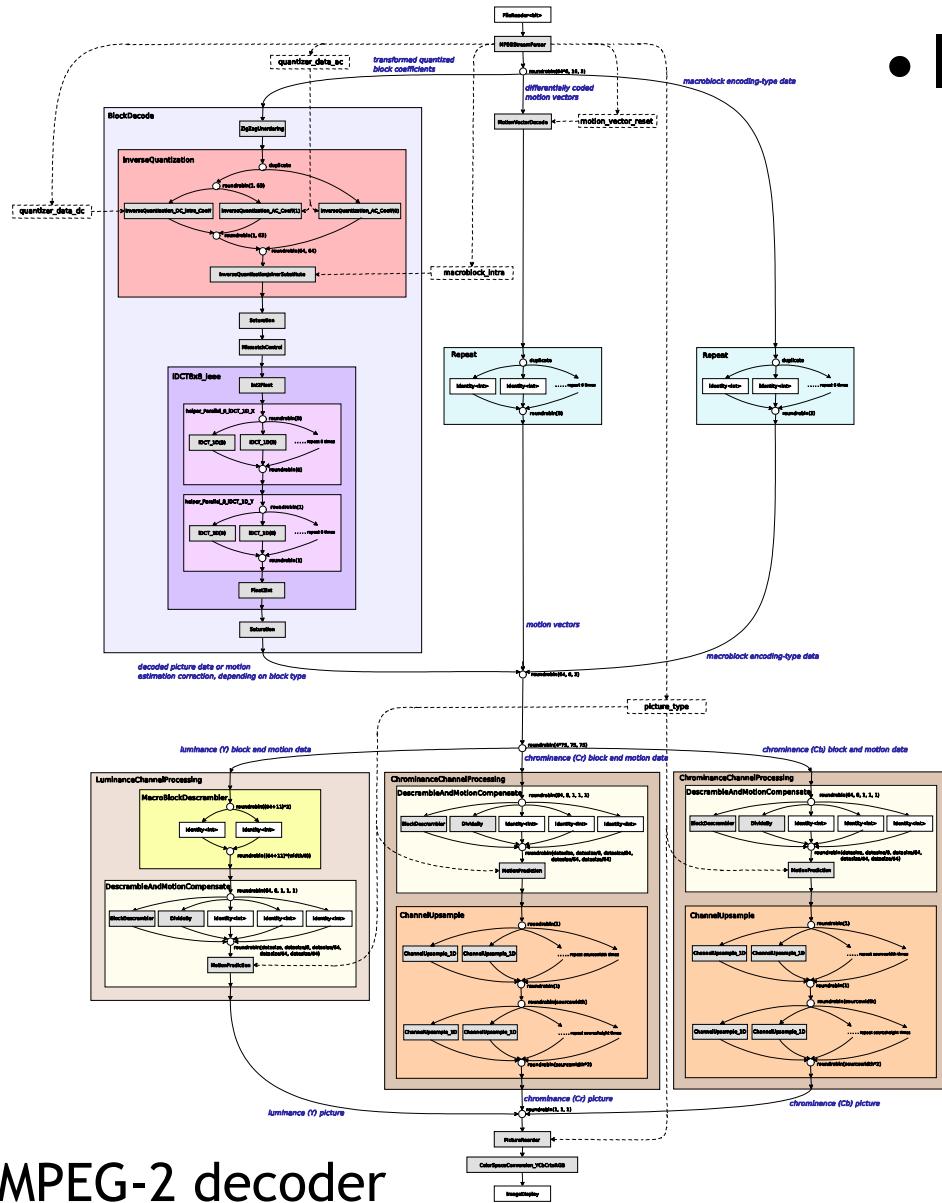


MPEG-2 decoder



DCT and quantization of 8x8 image block

# MPEG: Motion Video Codec



MPEG-2 decoder

- Implementation statistics
  - 4921 lines of code
    - 48 static streams
    - Compile to ~2150 filters
      - 352x240 resolution
  - Reference C implementation has 9832 lines of code
    - Supports interlacing and multi-layer streams
- 8 weeks of development
- 1 programmer with no prior MPEG-2 experience

# Excerpt from StreamIt Implementation

Specification in Section 7.4.1:  $F''[0][0] = \text{intra\_dc\_mult} \times QF[0][0]$

**Table 7-4 - Relation between intra\_dc\_precision and intra\_dc\_mult**

intra_dc_precision	bits_of_precision	intra_dc_mult
0	8	8
1	9	4
2	10	2
3	11	1

```
int->int filter InverseQuantization() {
    int[4] intra_dc_mult = {8, 4, 2, 1};
    int intra_dc_precision;

    work pop 1 push 1 {
        push(intra_dc_mult[intra_dc_precision] * pop());
    }
}
```

# Excerpt from Reference Implementation

Specification in Section 7.4.1:  $F''[0][0] = \text{intra\_dc\_mult} \times QF[0][0]$

**Table 7-4 - Relation between intra\_dc\_precision and intra\_dc\_mult**

intra_dc_precision	bits_of_precision	intra_dc_mult
0	8	8
1	9	4
2	10	2
3	11	1

```

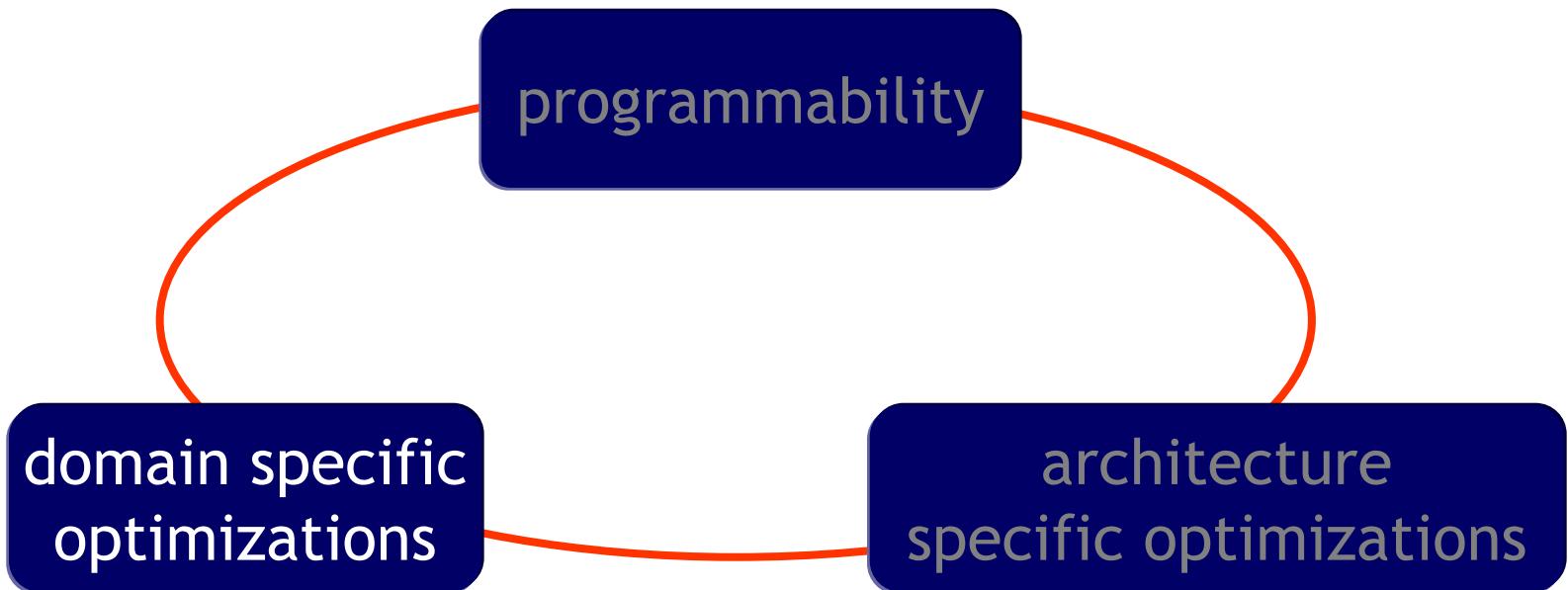
int[4] intra_dc_mult = {8, 4, 2, 1};

for (int m = 0; m < W*H/(16*16); m++)
    // six components for chrominance and luminance
    for (int comp = 0; comp < 6; comp++)
        if (macroblock[m].intra)
            macroblock[m].block[comp][0] *= intra_dc_mult[intra_dc_precision];
    // and many lines later
    if (cc == 0)
        val = (dc_dct_pred[0] += Get_Luma_DC_dct_diff());
    else if (cc == 1)
        val = (dc_dct_pred[1] += Get_Chroma_DC_dct_diff());
    else
        val = (dc_dct_pred[2] += Get_Chroma_DC_dct_diff());
    if (Fault_Flag) return;
    bp[0] = val << (3-intra_dc_precision);

```

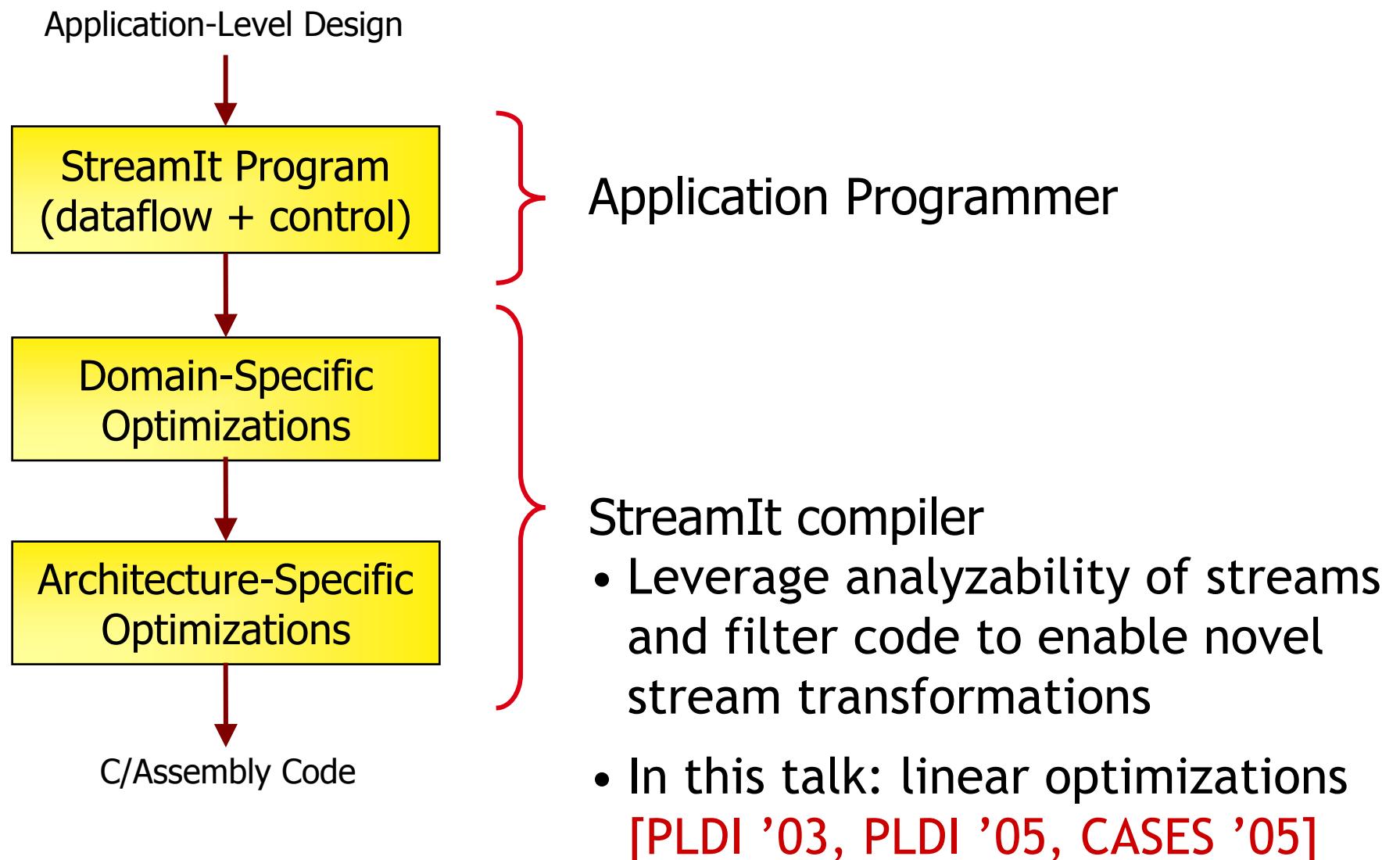
# The StreamIt Vision

- Boost productivity, enable faster development and rapid prototyping



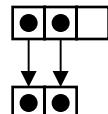
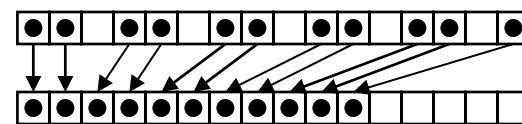
- Simple and effective optimizations for streams
- Targeting tiled architectures, clusters of workstations, DSPs, and traditional uniprocessors

# Design Flow with StreamIt

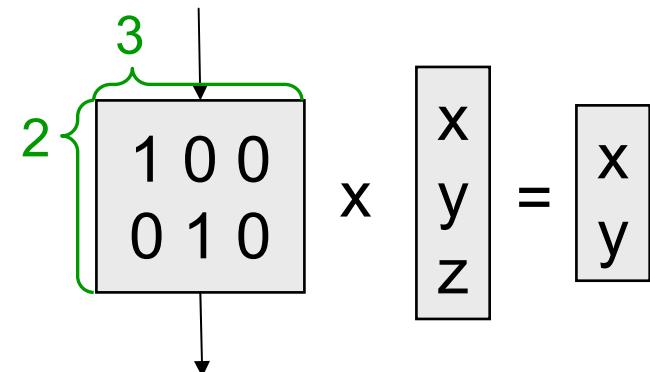


# Linear Filter Example

- “Drop every third bit in the bit stream”

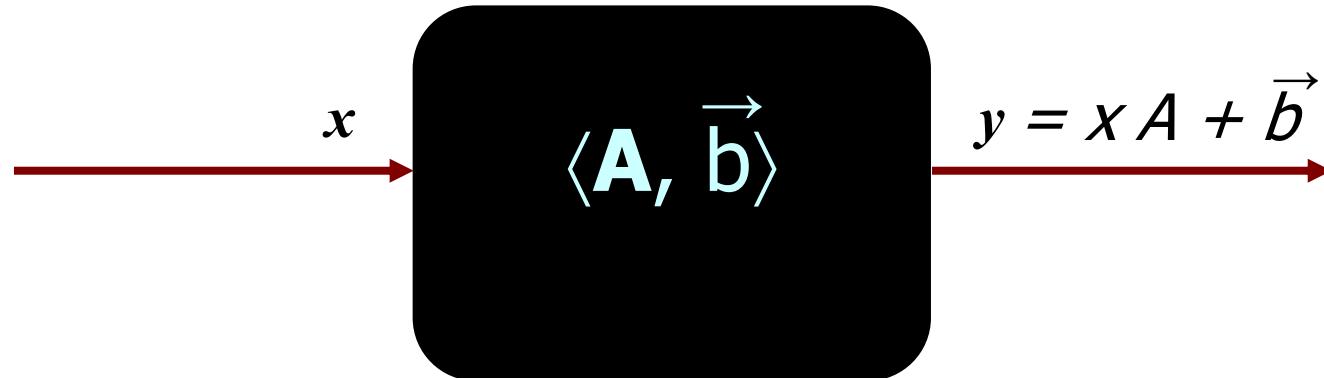


```
bit → bit filter DropThirdBit {
    work push 2 pop 3 {
        push(pop());
        push(pop());
        pop();
    }
}
```



# In General

- A linear filter is a tuple  $\langle \mathbf{A}, \vec{\mathbf{b}} \rangle$ 
  - $\mathbf{A}$ : matrix of coefficients
  - $\vec{\mathbf{b}}$ : vector of constants
- Example

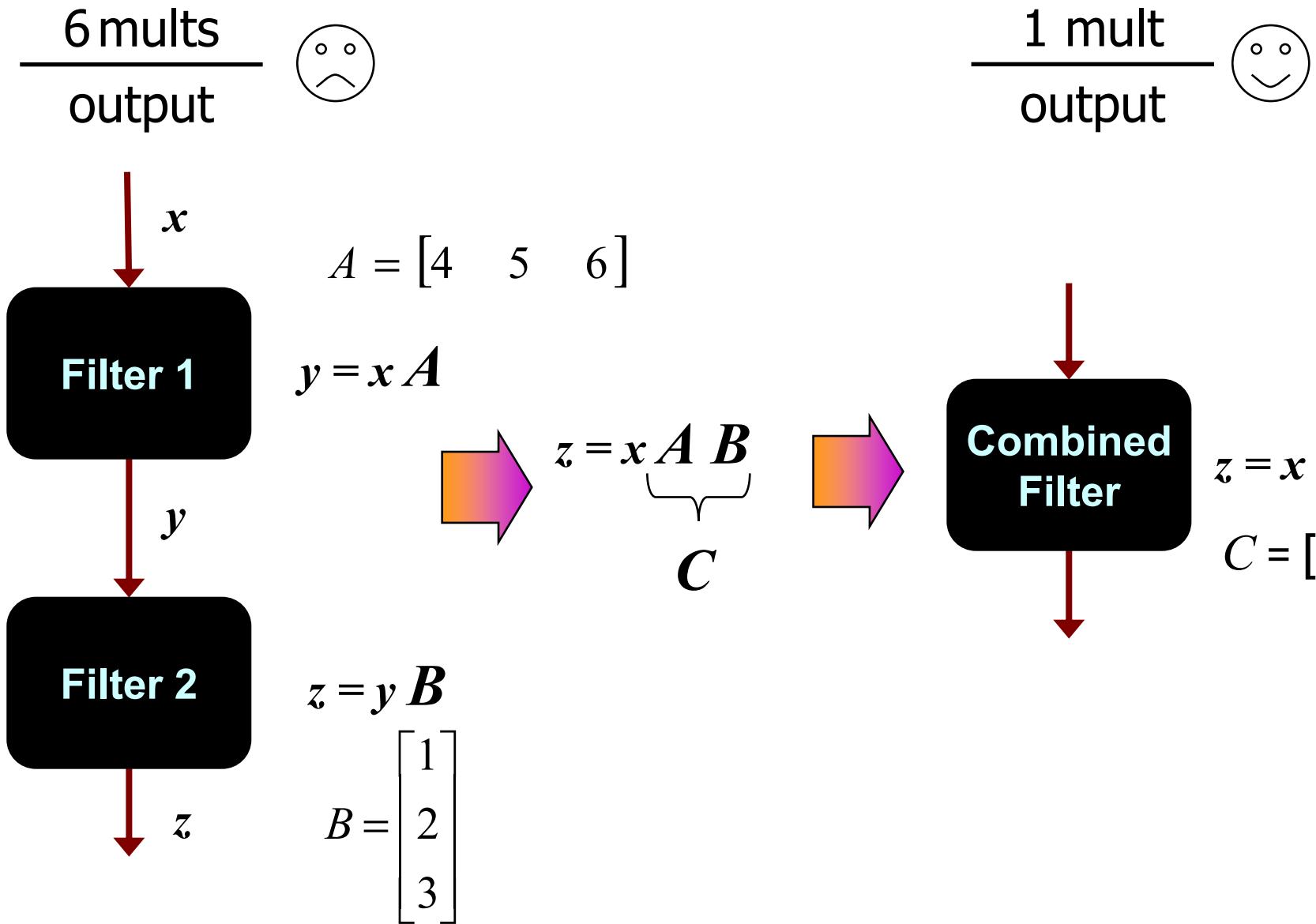


- Linear dataflow analysis resembles constant propagation

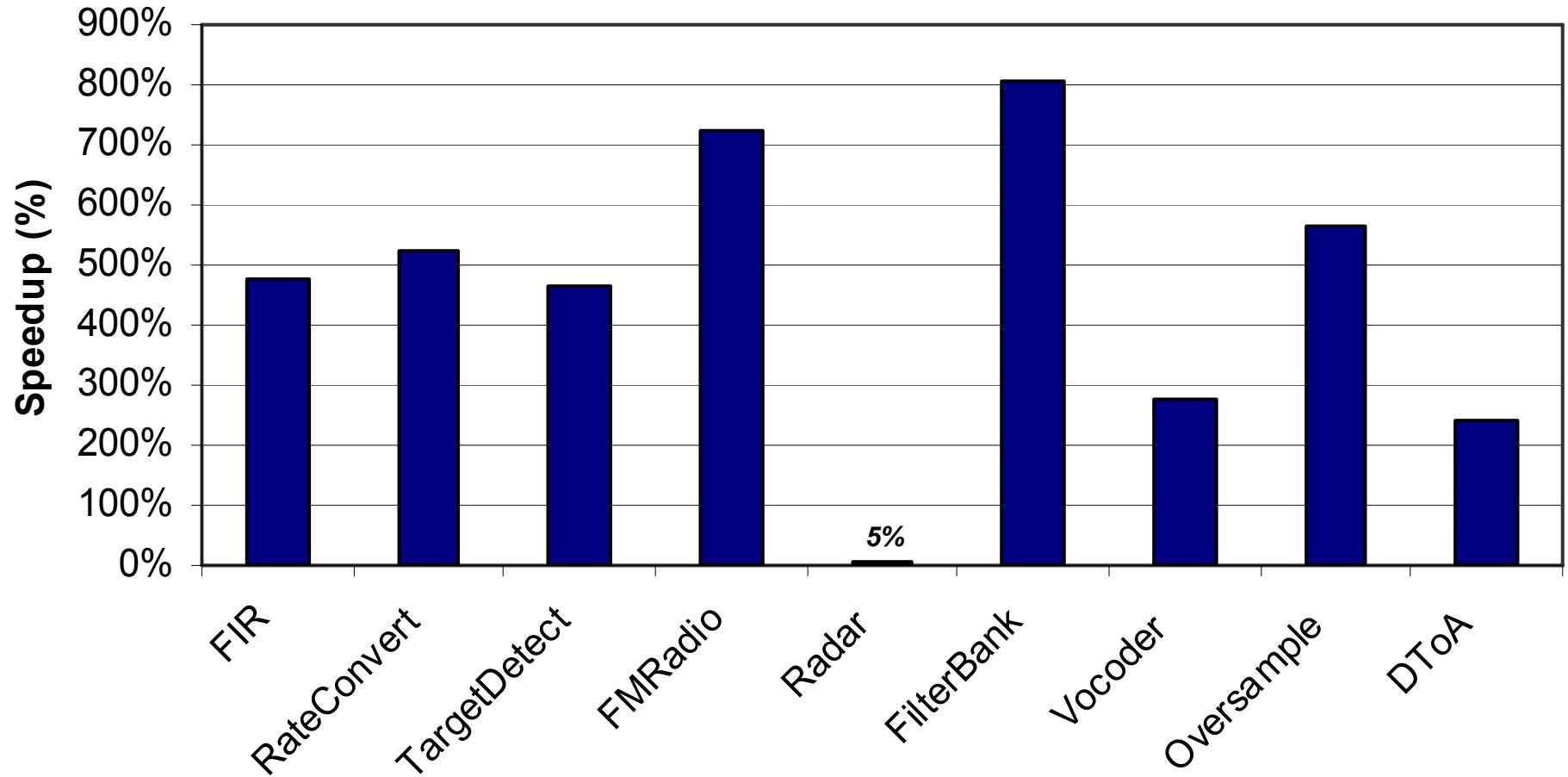
# Opportunities for Linear Optimizations

- Occur frequently in streaming codes
  - FIR filters
  - Compressors
  - Expanders
  - DFT/DCT
  - Bit permutations in encryption algorithms
  - JPEG and MPEG codecs
  - ...
- Example optimizations
  - Combining adjacent nodes
  - Also, translating to frequency domain when profitable

# Combining Linear Filters



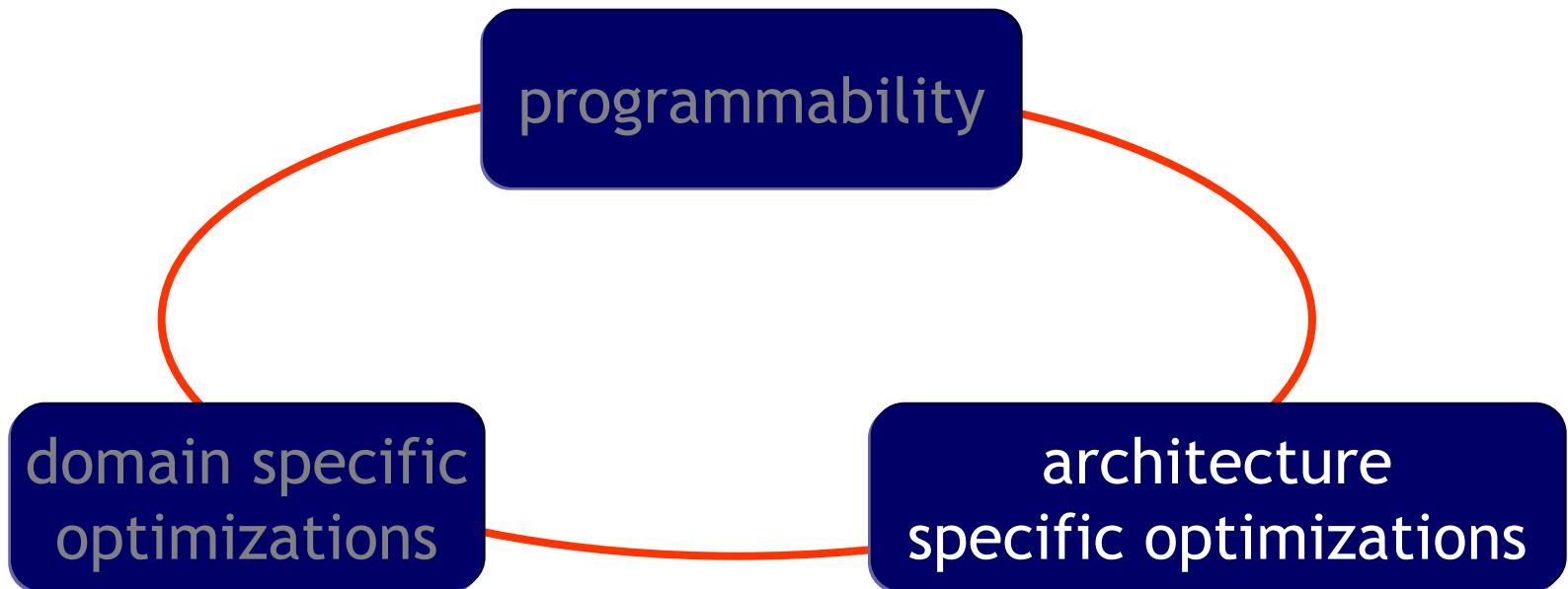
# Results from Linear Optimizations



Pentium 4 results compared to baseline StreamIt

# The StreamIt Vision

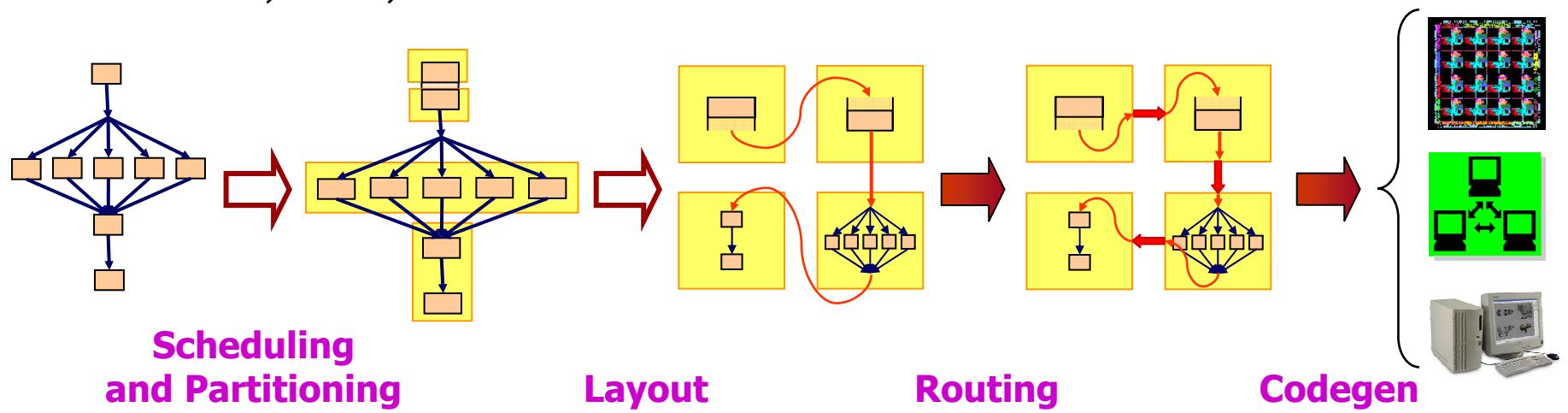
- Boost productivity, enable faster development and rapid prototyping



- Simple and effective optimizations for streams
- Targeting tiled architectures, clusters of workstations, DSPs, and traditional uniprocessors

# Core Compilation Technology

- Focused on a common challenges in modern and future architectures
  - MIT Raw fabric architecture
  - Clusters of workstations
  - ARM, x86, and IA-64



- Compiler's role: map the computation and communication pattern to processors, memories, and communication substrates

# Compiler Issues

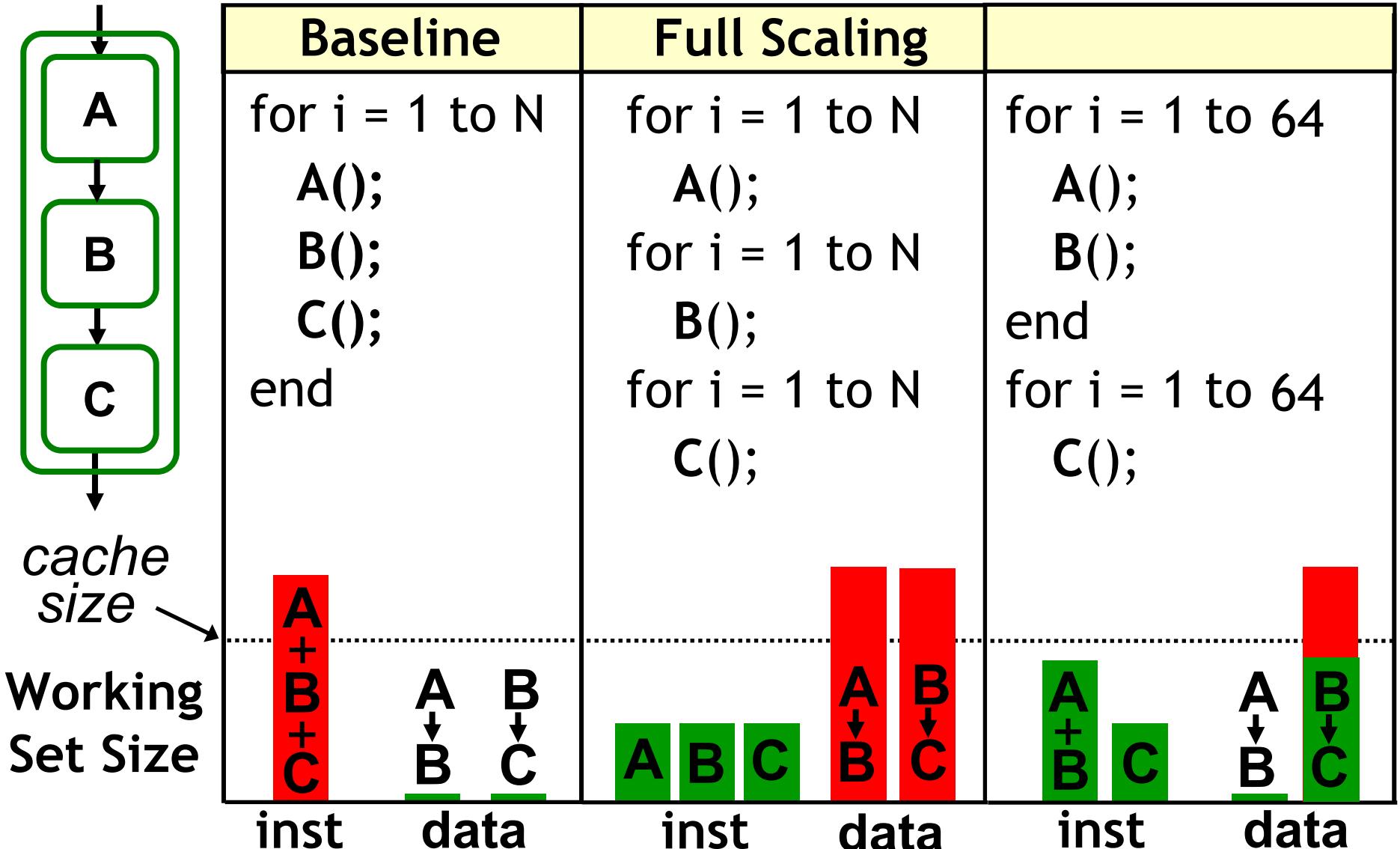
- Load balancing
- Resource utilization
- Fault tolerance
- Dynamic reconfiguration
- ...
- In this talk: cache aware scheduling and partitioning [LCTES '05]

# Example Cache Optimization

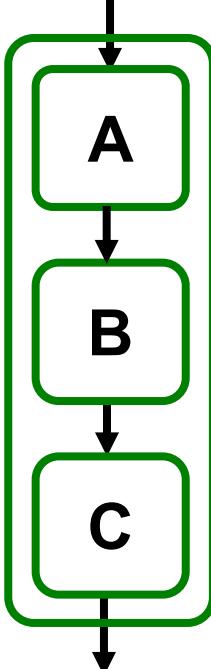
**Working Set Size**

	inst	data	inst	data	inst	data
cache size	A + B + C	A → B → C	A	B	A + B	A → B → C
Working Set Size						

# Example Cache Optimization



# Example Cache Optimization



*cache size* →

	Baseline	Full Scaling	Cache Aware
Working Set Size	<pre> for i = 1 to N     A();     B();     C(); end </pre>	<pre> for i = 1 to N     A();     for i = 1 to N         B();         for i = 1 to N             C();     end end </pre>	<pre> for i = 1 to 64 ;4     A();     B(); end for i = 1 to 64     C(); end </pre>

inst      data      inst      data      inst      data

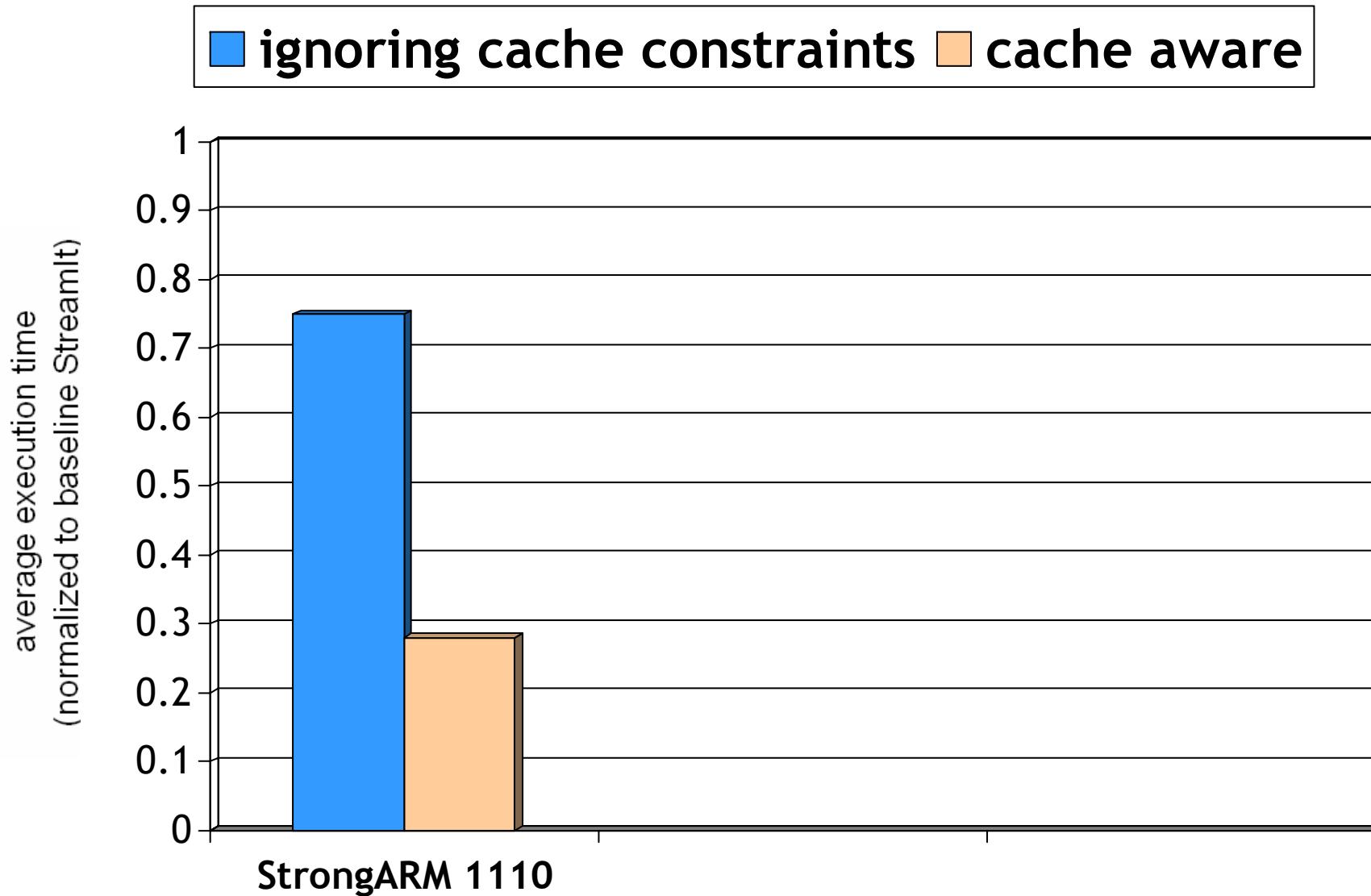
The diagram illustrates the cache behavior for three different scaling approaches. It shows the working set size (the number of items loaded into memory) for both instructions (inst) and data (data) across four stages.

- Baseline:** The working set for instructions is **A + B + C**. The working set for data is **A → B → C**.
- Full Scaling:** The working set for instructions is **A**. The working set for data is **B**.
- Cache Aware:** The working set for instructions is **A + B**. The working set for data is **C**.

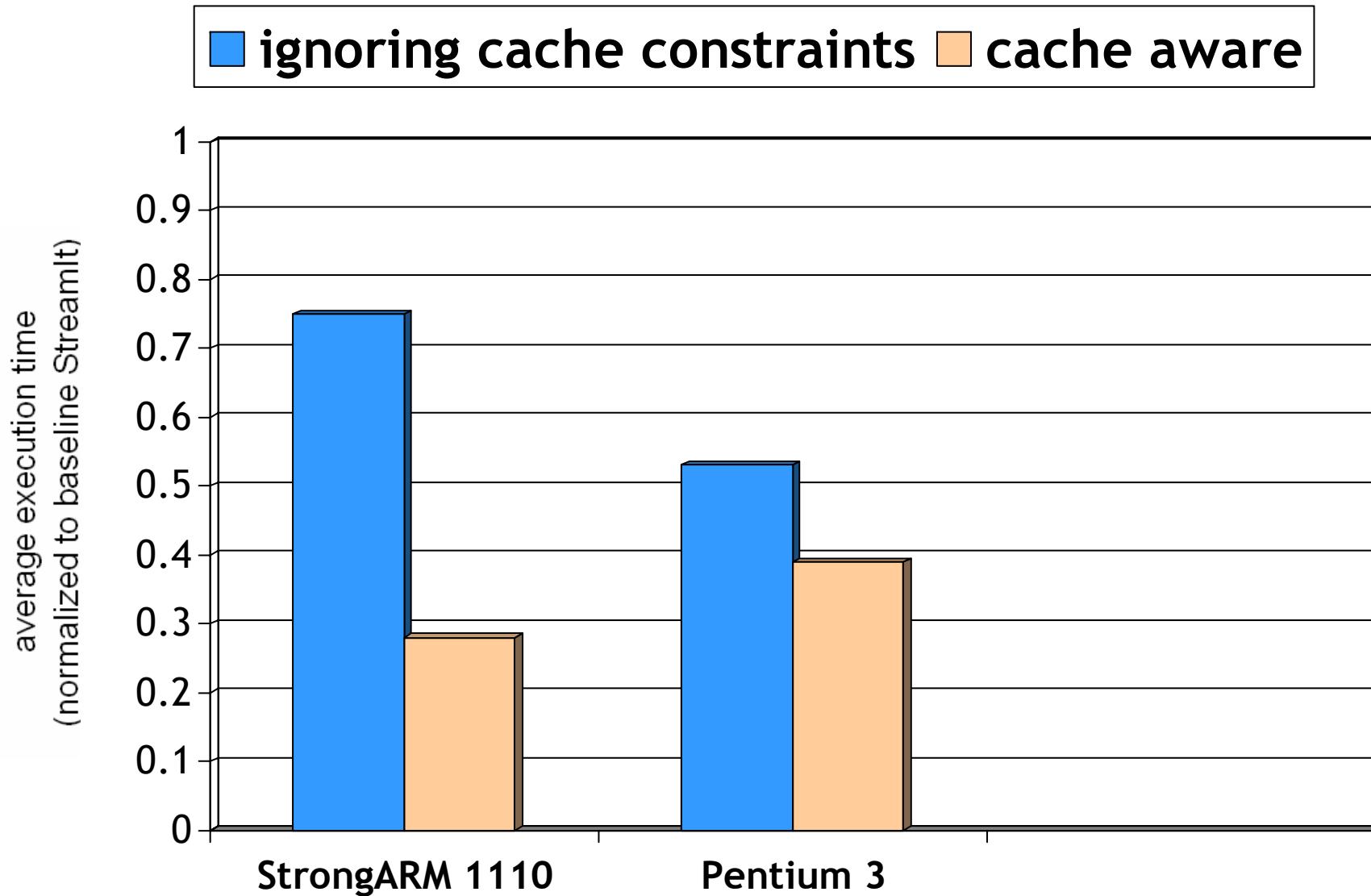
# Evaluation Methodology

- StreamIt compiler generates C code
  - Baseline StreamIt optimizations
    - Unrolling, constant propagation
  - Compile C code with gcc-v3.4 with -O3 optimizations
- StrongARM 1110 (XScale) embedded processor
  - 370MHz, 16Kb I-Cache, 8Kb D-Cache
  - No L2 Cache (memory 100× slower than cache)
  - Median user time
- Also Pentium 3 and Itanium 2 processors
- Suite of 11 StreamIt Benchmarks

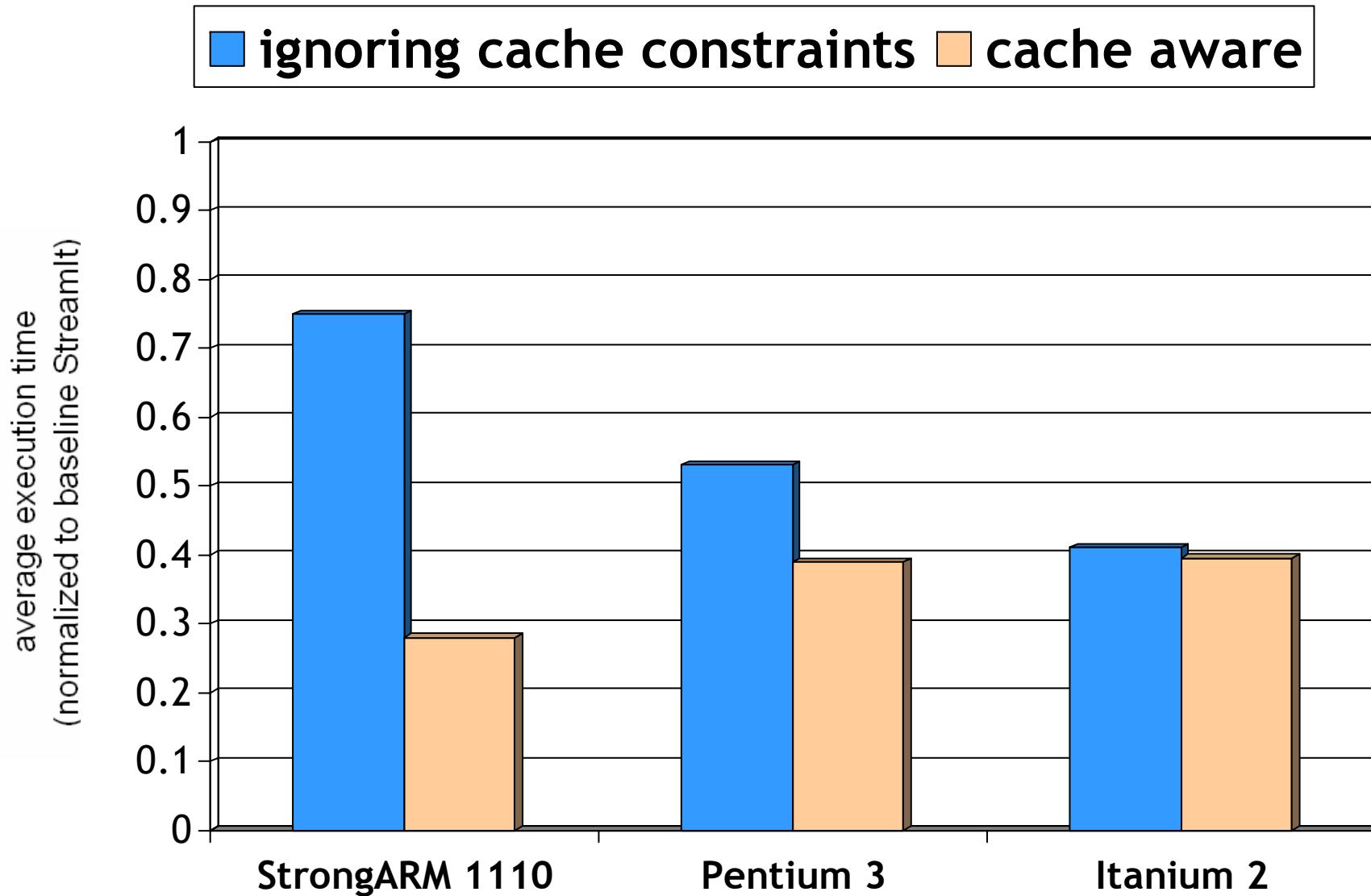
# Cache Optimizations Results



# Cache Optimizations Results

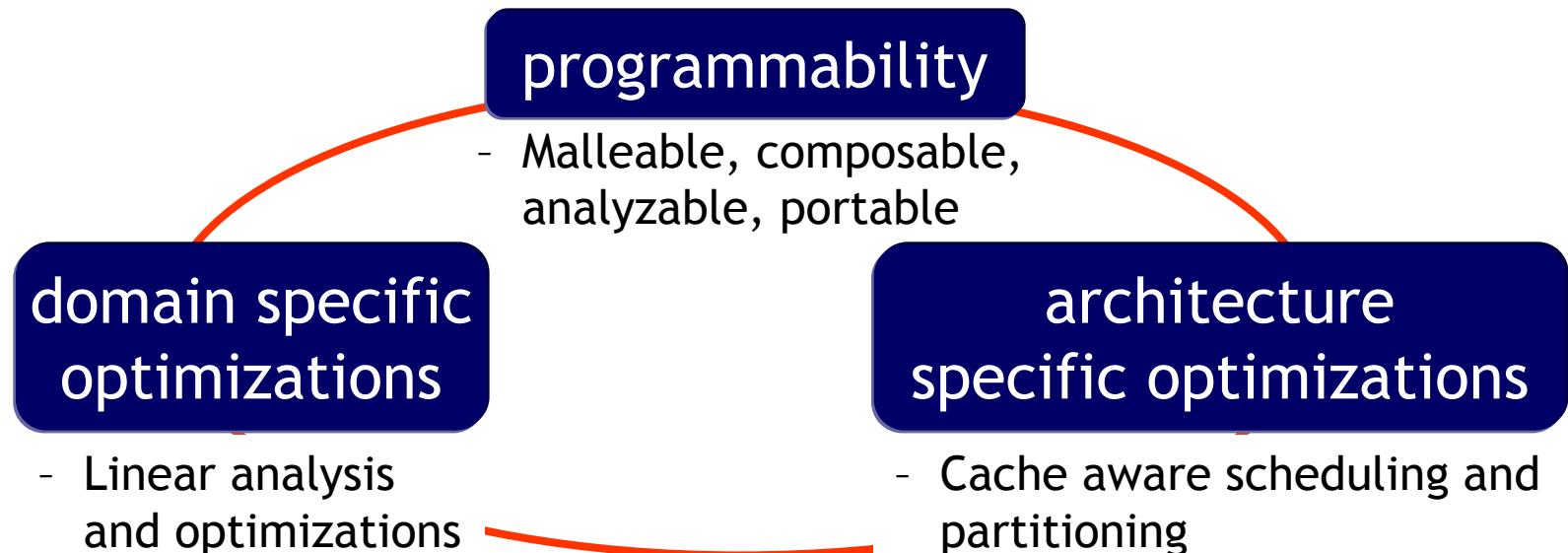


# Cache Optimizations Results



# Concluding Remarks

- StreamIt improves programmer productivity without compromising performance
  - Easily identify pipeline and data parallelism
  - Expose information for domain specific and architecture specific optimizations



# Broader Impact

- Integration into future HPCS languages
  - IBM: X10
- StreamIt for graphics applications
  - Programmable graphics pipeline [Graphics Hardware '05]
- StreamIt for emerging architectures
- Looking for users with interesting applications

# High-Productivity Stream Programming for High-Performance Systems

Rodric Rabbah, Bill Thies, Michael Gordon, Janis Sermulins,  
and Saman Amarasinghe

Massachusetts Institute of Technology



HPEC 2005, MIT LL