



Gedae: Auto Coding to a Virtual Machine

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What is Gedae?



Gedae is a block diagram language ...



Express signal and data processing algorithms, parallelism, load balancing, fault tolerance and mode control



..that Gedae transforms under user control...



User can set optimization parameters that are independent of the graph to guide transformation

Gedae ...to operate efficiently on a virtual machine.



Complete systems can be developed independent of the target system without losing runtime efficiency

Gedae Language



- Gedae provides application information through
 - modules with well-defined behavior
 - ports with well-defined characteristics
 - and manifest connectivity with explicit sequential and parallel execution paths
- This information is implicit in most languages
- Gedae makes the information explicit
 - over 50 different information expression features



Information provided by language allows Gedae to analyze and efficiently implement algorithms



Gedae Transformations

- The block diagram is transformed • using over 100 algorithms.
- The transformations establish the: •
 - Order of execution
 - Queue sizes
 - Granularities
 - Memory layout
 - Dynamic schedule parameters
 - Data transfer types and parameters
 - Mode control

Functional Implementation **Specification Specification Transformations Detailed Model** User Gedae Generation Vendor **Deployable** Application **Virtual Machine Heterogeneous HW**

The Gedae transformations build a detailed model of the deployed application. Gedae uses that information to provide visibility

Key





Gedae Virtual Machine (VM)

- Gedae provides the following components:
 - Command handler
 - Dynamic scheduler
 - Segmentation Support
 - Primitive Support
 - Visibility Support
- The vendor provides
 - Inter-processor communications
 - Optimized vector libraries
 - Other basic services



The Gedae virtual machine makes applications processor independent

Three Examples



- Real-Time Space-Time Adaptive Processing (RT-STAP)
 - Miter benchmark graph
 - Illustrates efficient parallel execution of large graph
- Multilevel Mode Graph
 - Illustrates nested mode control with distributed state
 - Dynamic data application
- Sonar Graph
 - Illustrates large data reduction during processing

Each example illustrates features of the language, transformations, and virtual machine



RT-STAP: Language



Families permit replicating box and data elements



RT-STAP: Language



- Instantiation constants control the size of the graph
- Routing boxes allow equation based connectivity



RT-STAP: Transformations

- User maps primitives to physical processors
- Gedae transforms graph by inserting send/receive primitives to communicate between partitions
- Gedae automatically creates executables to run on each processor

🚊 Group 1 Partition Table									
Fi	e <u>E</u> dit <u>∨</u> iew	Options							
Nar	ne	Part	SubSched						
	[55]applyW	113							
	[56]арр1уЫ	114							
	[57]applyW	114							
	[58]app1yW	114							
	[59]app1yW	114							
	[60]app1yW	115							
	[61]applyW	115		٦I					
	[62]app1yW	115							
	[63]app1yW	115							
	[0]form_W	100							
	[1]form_W	100							
	[2]form_W	100							
	[3]form_W	100							
	[4]form_W	101							
	[5]form_W	101							
	[6]form W	101							

Different mappings can be tried without modifying the graph – the needed transformation happens automatically

RT-STAP Transformations



- User can set transfer properties on send/recv pairs with Transfer Table
- Transformations automatically set parameters to send/recv pairs to communicate these properties to running application

Group 1 Transfer Table								_ 🗆 🗙
<u>File Edit View Options</u>								Help
Name	Source	Dest	Xfer Type	NBsize	Send Bufs	Recv Bufs	Xfer Params	
stap_pp.form_Xp.corner_t.[0][13]mx_vx_1 <in< td=""><td>116</td><td>113</td><td>stream</td><td>46080</td><td></td><td></td><td>-b 10 -d 2</td><td></td></in<>	116	113	stream	46080			-b 10 -d 2	
stap_pp.form_Xp.corner_t.[7][14]mx_vx_1 <in< td=""><td>123</td><td>114</td><td>dsa</td><td></td><td>2</td><td>3</td><td>-b 10 -d 2</td><td></td></in<>	123	114	dsa		2	3	-b 10 -d 2	
stap_pp.form_Xp.corner_t.[6][14]mx_vx_1 <in< td=""><td>122</td><td>114</td><td>dsa</td><td>8</td><td>2</td><td>3</td><td>-b 10 -d 2</td><td></td></in<>	122	114	dsa	8	2	3	-b 10 -d 2	
stap_pp.form_Xp.corner_t.[5][14]mx_vx_1 <in< td=""><td>121</td><td>114</td><td>dsa</td><td></td><td>2</td><td>3</td><td>-b 10 -d 2</td><td></td></in<>	121	114	dsa		2	3	-b 10 -d 2	
stap_pp.form_Xp.corner_t.[4][14]mx_vx_1 <in< td=""><td>120</td><td>114</td><td>dsa</td><td></td><td>1</td><td>1</td><td></td><td>12</td></in<>	120	114	dsa		1	1		12
stap_pp.form_Xp.corner_t.[3][14]mx_vx_1 <in< td=""><td>119</td><td>114</td><td>stream</td><td>46080</td><td></td><td></td><td></td><td></td></in<>	119	114	stream	46080				
stap_pp.form_Xp.corner_t.[1][15]mx_vx_1 <in< td=""><td>117</td><td>115</td><td>stream</td><td>46080</td><td></td><td></td><td></td><td></td></in<>	117	115	stream	46080				
stap_pp.form_Xp.corner_t.[0][15]mx_vx_1 <in< td=""><td>116</td><td>115</td><td>stream</td><td>46080</td><td></td><td></td><td></td><td></td></in<>	116	115	stream	46080				
stap_pp.form_Xp.[0]rfam_mx<[0]in	115	100	stream	84480				
stap_pp.preprocess2.[20]iq_convert2.v_multOsc <in< td=""><td>124</td><td>123</td><td>dsa</td><td>2</td><td>1</td><td>1</td><td></td><td></td></in<>	124	123	dsa	2	1	1		
stap_pp.preprocess2.[21]iq_convert2.v_multOsc <in< td=""><td>124</td><td>123</td><td>dsa</td><td>6</td><td>1</td><td>1</td><td></td><td></td></in<>	124	123	dsa	6	1	1		
sink_pp.[0]vx_mux<[0]in	100	host	gsim_host>host	3840				
sink pp.[0]vx mux<[1]in	100	host	asim host≻host	3840		1		N.

User can guide transformations to optimize implementation

RT-STAP: Running on VM



rt_stap_hard_8_16 Trace Ta	ble		
<u>File View Options</u> St	ats 🔺 🗾		
Name	9,226810 s	385269 s	10.034096 s +0.498933 s
client			
100		1-1	
101			
102			
103			
104			
105			
106			
107			
host			
	M	4	

Send/Recv webs show interprocessor communication and uncover synchronization problems

RT-STAP: Running on VM



Preplanned use of memory allows distributed runtime debugging

Mode Control: Language



- "Exclusive" branch outputs show where resources can be shared
- State shared between modes is explicitly declared in the graph

The Gedae primitive language directly supports segmented data processing, sharing of resources, and distribution of state



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Mode Control: Language

- Branch box copies input data stream to one of a family of outputs based on a control stream. Output is:
- Segmented the box will add segment boundaries to the output
- **Dynamic** the box will state how much data is **produced** on the output at runtime.
- Exclusive only one of the family of F outputs gets data on any firing. Allows sharing of resources and state.

The Gedae extensible language has no "built-in" primitives. 8000+ delivered primitives. Users can add custom primitives

```
Name: cp branchf e
             Input:
                      stream ControlParamRec in;
             Input:
                      stream int c;
             Tocal: int last;
             Output: exclusive segmented dynamic stream
                      ControlParamRec [Flout;
             Reset: { last = -1; }
             Apply: {
                int q,i;
                int prdc = 0;
                for (q=0; q<granularity; q++) {</pre>
                  int j = c[q];
                  if (last != j) {
                    if (0 \le \text{last} \& \text{last} \le F) {
                      produce(out[last],prdc);
                      prdc = 0;
                      segment(out[last],SEGMENT END);
                    last = i;
                  if (0<=j && j<F) {
                    *out + + = *in;
                    prdc++;
                  in++;
                produce(out[last],prdc);
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```

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Mode Graph: Transformation



User can set partitioning, mapping, data transfer methods, granularity, priority, queue sizes and schedule properties from the group control dialog

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Mode Graph: Running on VM



• Each mode requires a different number of processors

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• Branch boxes at one level are responsible for the dynamic distribution

VM runtime kernel enforces dynamic data driven execution. Send and receive primitives and state transfer primitives use BSP of virtual machine to transfer data

Mode Graph: Running on VM



- Primitives to send and receive state are automatically added by transformations
- Messages generated by Virtual Machine at mode change boundaries efficiently coordinate state transfers



Result is efficient transparent use of shared state on distributed processing system

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Sonar: Language



Sonar Graph creates low bandwidth output from high bandwidth input data

Sonar: Language



- Connectivity + Port Descriptions gives information needed to schedule graph
- mx_vx produces R=120 tokens out for every 1 token in
- vx_multV box must fire 120 times for each firing of the mx_vx box.
- vx_fft box fires one time for each firing of vx_multV box
- Simple predetermined schedule generated from graph and info embedded in primitives





Can create a multirate graph that has boxes firing at different granularities Gedae, Inc. www.gedae.com

Sonar: Transformation



- User can place boxes in subschedules to strip-mine the vector processing
- Allows use of fast memory
- Can reduce memory usage



Multirate graphs can be implemented using subscheduling to improve speed and reduce memory usage Gedae, Inc. www.gedae.com

Sonar: Transformation

Auto-Subscheduling Tool

🔀 Gra	oup 1 G	ain Hier	Table							_ □	1 ×
File	<u>E</u> dit	View	Options							Н	elp
Name			TotalG	Bytes	Div	TDiv	G	Bytes*G	Boxes	Subsched	
Sched	ule 1		1	1440	1	1	1	1440	6		
1	42		2	1966080	2	2	1	1966080	2	1	
	2		240	20484	30	60	4	81936	6	1.1	
	3		1024	4800	32	64	16	76800	1	1.2	
	4		4096	63628	64	4096	1	63628	17	1.2.1	Grou

- User can put boxes into named subschedules manually – but can be difficult
- Auto-Subscheduling Tool puts boxes in subschedules automatically
- Finds nested sets of connected boxes running at common granularities.
- Automatically sets subscheduling levels

Auto-subscheduling has reduced memory needed by graph from 250 Mbytes to about 2.5 Mbytes - 100x improvement Gedae, Inc. www.gedae.com



Schedule	Information
Dialog	

Schedule Parameters

File	Ē	dit	\geq	iew	<u>O</u> ptions					
Name						Size	Priority	Policy	Peri	
Part default										
Sch	Schedule 1				1448	0	dataflow			
	SubSched 1									
		Seg	gment default			2457600				
		Seg	gment parent_memory				480			
		Sut	bSched 1.1							
			Seg	gment	default		65536			
			Segment parent_memory			ory	32784			
		Sut	ibSched 1.2							
			Seg	gment	default		143040			
			Segment parent_memory			ory	15360			
			Sut	Sched	1.2.1					
				Segme	nt default		61720			
				Seqme	nt parent r	nemoru	960			

Sonar: Running on VM

Gedae

test_WB_SONAR Trace Table		
File View Options Stats 🖌 🗾		
Name	43.247102 s 43.184982 s	43,696022 s +0,358640 s
Client		
default		
Schedule 1		
SubSched 1		
SubSched 1.2	I	
BasebandFilterDec		
td_beamformer		
[0]v_spaceTimeSto	re_n 	
[0]interpolation		
[0]m_cselem		
v_nconcat		
WB_SONAR		
BasebandFilterDec		
FreqTransformation		
SubSched 1.1	-	
FregTransformation		
MakeWB		
MakeWB		
WB_SONAR		
DispWFandPlotwithControlS		
1		M

Multiple levels of subscheduling evident on Trace Table

Conclusion



- Gedae Block Diagram Language allows simple expression of a wide range of algorithms
- User optimization information can be added without modifying block diagram
- 100+ transformations create efficient executable application from language and user information
- Application runs efficiently on Virtual Machine
- VM provides portability and visibility