

# The *Mercury* System: Embedding Computation into Disk Drives

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## 1. Introduction

Having inexpensive data storage has enabled the amassing of vast amounts of information. At present, these data sets far exceed the capacity of modern processors, so searching them has become a serious challenge. In a recent invited talk at the High Performance Embedded Computing Workshop, John Reynders of Celera Genomics commented that, “The size of the databases we deal with is no longer measured in terabytes, but in exabytes.” [1]

The *Mercury* system is a prototype data search engine that can be embedded within the disk drive itself. We focus on the specific problems associated with searches of unstructured, unindexed data. Three specific applications include approximate matching of text (important for text searches of specific interest to homeland security where the original alphabet is different than the Latin alphabet and transliteration is involved), genomics and proteomics searches (important biological applications), and image searches (also of significant interest for homeland security).

Currently, data searching applications are implemented using traditional, off-the-shelf hardware platforms. Figure 1 illustrates the relevant features of virtually all of these systems. A disk (actually many disks) is attached via a controller to the I/O bus of a computer system. A path (labeled “bridge” in the figure) exists that enables data to flow from the I/O bus to the memory bus (and therefore to the system’s main memory). When an algorithm is executed on the processor, references to the main memory cause the data to be loaded into cache, at which point the processor can efficiently access the data.

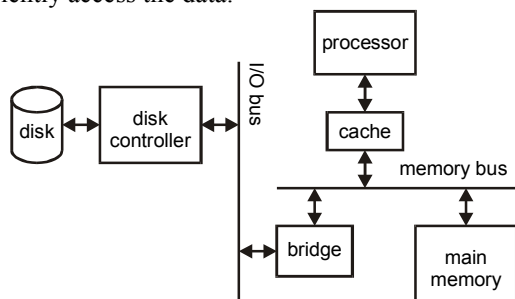


Figure 1. Typical hardware platform for data searching applications.

Add the operating system overhead to the above data movement requirements, and it is clear that there are significant data movement inefficiencies in current systems. Yet this is the system environment associated with development and deployment of virtually all of today’s data search applications. The result is that even though individual components in the system are quite fast (e.g., modern processors have clock speeds exceeding 2 GHz), the overall performance suffers because these fast components are used inefficiently.

Our system dramatically increases the speed with which large volumes of data can be searched by eliminating the above inefficiencies and searching data much closer to where it resides, on the disk, and by performing low-level search operations directly in reconfigurable hardware. In effect, the computation is being embedded into the drive.

## 2. Overall System Architecture

The *Mercury* system architecture is illustrated in Figure 2. Associated directly with a disk head is a Data Shift Register (DSR) that receives data streaming off the head at disk rotational speeds. The data in the DSR is made available (in parallel form) to reconfigurable logic that performs low-level searching operations on the data that has been retrieved off the disk. The specific function performed by the reconfigurable logic is tailored to the particular application of interest. Example functions for a set of applications are described below.

Also present in the system is a microprocessor that is used both for control duties (e.g., managing the function of the reconfigurable logic) and higher-level data searching operations. The remainder of the system reflects traditional designs, with an I/O bus, a bridge to the memory bus, and a classic memory hierarchy. The host processor is still responsible for managing the file system and maintaining the general functionality of the database, the new embedded hardware is used primarily for high-volume data searching operations.

In the description that follows, we assume (for discussion purposes) that the data is unstructured text and we are performing string matching queries. Realistic queries are generally compound in nature. In this case the search involves both matching the query strings to documents on the disk, and determining if the relationships exist when matches occur (or don’t occur).

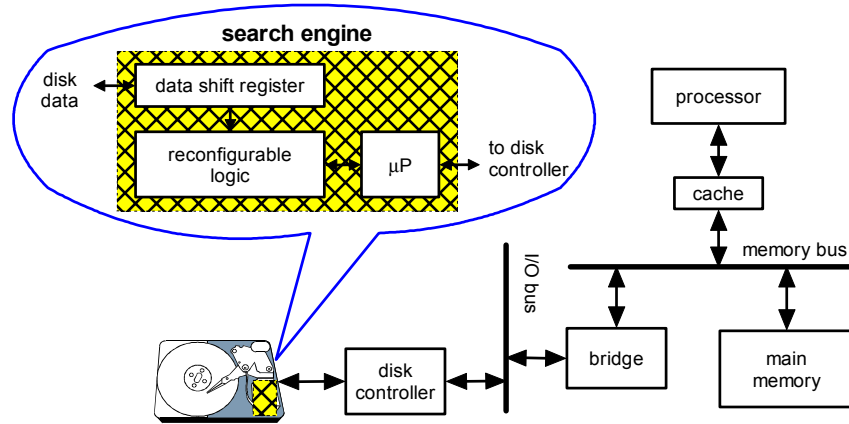


Figure 2. Overall system architecture.

For example, a query might contain the strings “Iraq”, “Iran”, “Israel” and “France”, and want to identify documents where either “Iraq” *and* “Israel”, *or* “Iran” *and* “France” are present. This form of compound inquiry may result in a number of document matches, and the objective is to perform these matches at disk speeds. In general, such queries can be represented as a tree structure where the leaves of the tree correspond to the byte strings being sought (e.g., Iraq, Israel, etc.), and the nodes correspond to the logical operations required of the query. One way of viewing the processing of a compound query is in terms of two components:

- processing of the leaves associated with the bottom of the tree (i.e., taking the leaf words and performing a match with data on the disk), and
- performing the logic operations associated with the combining nodes in the query tree structure.

In this architecture, leaf processing (word matching) is done in the reconfigurable hardware (at disk speeds) and the results are sent to the dedicated control microprocessor that acts to execute the logic associated with the nodes of the query tree. The result of this logic execution is a set of results that are sent to the host processor that initiated system activity by sending the original query.

### 3. Example Applications

To illustrate the use of the system, a number of example applications are described: unstructured text searches, biological sequence matching, and image searches. All of these applications are currently heavily used, and all tax the capabilities of current systems to deliver the throughput desired.

**Unstructured text searching:** As described above, a compound query posed in a text search context can be decomposed into the individual word searches (to be executed in the reconfigurable logic) and the composition of the word search results (executed on the

microprocessor). When requiring an exact match at the word level, it is sufficient for the reconfigurable logic to contain a register that stores the string to be matched and a comparator that indicates when the contents of the data shift register equal the string to be matched. We have an implementation of the above design operating with a data throughput of 4 Gb/s (i.e., the system can accept data at that rate sustained).

As an alternative to exact string matches, it is often useful to express the desired search in terms of an acceptable number of mismatches (e.g., insertions, deletions, and substitutions). Algorithms for this type of search (incorporated into the Unix command `grep`) are described in [2]. Our systolic array implementation of this algorithm executes at a clock rate of 100 MHz, accepting one character each clock, for an aggregate data rate of 800 Mb/s.

**Sequence matching:** The basic set of operations in genomic or proteomic sequence matching corresponds to a dynamic programming problem when executed on a conventional system [3]. Here, a pattern  $p$  is compared against symbols from a target  $t$ , and  $d_{ij}$  (an entry in the dynamic programming table) represents the edit distance at position  $i$  in the pattern,  $p_i$ , and position  $j$  in the target,  $t_j$ . In normal usage the pattern is short relative to the target. Typical sizes might have  $p$  on the order of 1000-2000 characters and  $t$  many billions of characters.

The details of our systolic array implementation of the dynamic programming problem are described in [4]. It can sustain a data throughput of 800 Mb/s.

**Image searching:** Many searching applications operate on data that represent a two-dimensional entity, such as an image. For example, one approach to the object recognition problem is to repeatedly compare the field of interest in an image to templates that store a representation of the objects to be recognized [5]. For imaging applications, the structure of the reconfigurable logic from Figure 2 must take into account the fact that the logical structure of the data is two-dimensional. In addition, the matching operations themselves are often

significantly different on two-dimensional data, and this must also be supported by the reconfigurable logic.

As an example of the type of image search problem that can be accelerated with the proposed system, we have investigated the automatic recognition of objects within synthetic radar (SAR) imagery [5,6]. Here, SAR images are compared with templates in an object database using a conditionally Gaussian data model. An implementation of the reconfigurable logic for this application is under development.

Another potential use of the image search system is object recognition in millimeter wave imagery of people entering a restricted area, such as an airport concourse. In these images, non-metallic sharp objects are readily detected, yet there are serious privacy concerns associated with human scanning, since the modality effectively “sees” through clothing. An automated recognition system that does not present images to operators has the potential to be effective at the security task while mitigating (though not eliminating) the privacy issue.

**Other applications:** The above section has described the reconfigurable hardware operation for a few specific applications. Other applications will have distinct low-level operations that need to be performed. One strength of this system is the fact that the reconfigurable logic is flexible enough to support not only the example applications already described, but a large variety of additional structures, even those not yet envisioned when the system was designed.

## 4. Performance

While the overall system is still under development, a sufficient number of component pieces exist to report on their performance. Three reconfigurable hardware search kernels have undergone detailed design (i.e., VHDL-level design), and all are operational. An exact text search engine design operates at 62.5 MHz (eight characters per clock, or 4 Gb/s); the approximate text search engine (agrep) has been tested to 100 MHz (one character per clock, 800 Mb/s); and the biosequence search engine has been tested to 25 MHz (four characters per clock, 800 Mb/s) [4].

Software implementations of these three kernels have been measured on a 933 MHz PC as follows: 280 Mb/s for exact text searches, 26 Mb/s for approximate text searches (allowing up to 8 errors), and 6.4 Mb/s for the biosequence search. Table 1 shows the speedup achievable under two conditions, one is the measured performance as limited by our current ATA (7200 rpm) drive and the other under the conditions where a faster drive is available and the reconfigurable logic kernel is the performance limit.

**Table 1. Application speedups.**

Application	Disk-limited speedup	Logic-limited speedup
Exact text search	1.1	14
Approx. text search	12	31
Biosequence search	50	125

## 5. Summary and Conclusions

This paper presents the basic design of a data-mining system that has the potential for truly fast operation, unhindered by the overheads imposed by the I/O bus, main memory bus, cache, operating system, etc. An important requirement for the ultimate success of this system is the decomposition of data search operations into low-level operations that can execute on the reconfigurable hardware and high-level operations that execute in software.

The system achieves performance gains via four mechanisms: reduced data movement overhead, searches operating at hardware speeds, specialization of the hardware logic to the particular query, and parallelism at both the disk and system levels.

We currently have initial prototype implementations of the reconfigurable hardware for several of the applications described above, and are working to further improve their performance. We are also developing performance models that help assess the performance gains that can be achieved using the system.

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- [3] Dan Gusfield, *Algorithms on Strings, Trees, and Sequences*, Cambridge University Press, 1997.
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